## Converting between binary and decimal

To easily convert a number represented in binary notation, such as $10100_{[2]}$, we can employ Horner's algorithm. At each step, we multiply the previous result by 2, and add the next bit in the number. To convert in the other direction, we divide by 2 and write the remainder at each step from bottom to top. We can see the conversion between $10100_{[2]}$ and 20 (or $20_{[10]}$ to be extra-decimaly) below.


## Checkpoint 0

What is the decimal representation of $1111010_{[2]}$ ?
What is the binary representation of $49_{[10]}$ ?

## Hexadecimal notation

Hex is useful because every hex digit corresponds to exactly 4 binary digits (bits). Base 8 (octal) is similarly useful: each octal digit corresponds to exactly 3 binary digits. However, hex more evenly divides up a 32 -bit integer. In C0 we indicate we are using base 16 with an $0 x$ prefix, so $7 f 2 c_{[16]}$ is $0 x 7 f 2 c$.

| Hex | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | a | b | c | d | e | f |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Bin. | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | 1000 | 1001 | 1010 | 1011 | 1100 | 1101 | 1110 | 1111 |
| Dec. | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |

Convert the binary number $1011111010101101_{[2]}$ to hex.
Convert the $0 \times 20$ to decimal.
Why wouldn't it make sense to write a C0 function that converts hex numbers to decimal numbers?

## Bit manipulation

| and |  |  | or |  |  | xor (exclusive or) |  |  | complement |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| \& | 1 | 0 | 1 | 1 | 0 | $\wedge$ | 1 | 0 | $\sim$ | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |  | 0 | 1 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 |  |  |  |

There are also shift operators. They take a number and shift it left (or right) by the specified number of bits. In C0, right shifts sign extend. This means that if the first digit was a 1 , then 1 s will be copied in as we shift.

$$
1101111101010010_{[2]} \gg 8=\mathbf{1 1 1 1} \mathbf{1 1 1 1} 11011111_{[2]}
$$

## Checkpoint 1

What does $\left(00010101_{[2]} \& 00110101_{[2]}\right) \mid\left(10101010_{[2]}^{\wedge} 00011110_{[2]}\right)$ evaluate to?

What does $\left(5_{[10]} \mid 13_{[10]}\right)^{\wedge}\left(28_{[10]} \& 10_{[10]}\right)$ evaluate to?
What is the difference between logical and bitwise operators?

## Two's complement

Because C0's int type only represents integers in the range $\left[-2^{31}, 2^{31}\right.$ ), addition and multiplication are defined in terms of modular arithmetic. As a result, adding two positive numbers may give you a negative number!

## Checkpoint 2

Write a function that returns 1 if the sign bit is 1 , and 0 otherwise. That is, write a function that returns the sign bit shifted to be the least significant bit. Your solution can use any of the bitwise operators, but will not need all of them.

```
int getSignBit(int x)
//@ensures \result == 0 || \result == 1;
```

3 \{
$\qquad$ ; 5 \}

## Checkpoint 3

What assertion would you need to write to ensure that an addition would give a result without overflowing (in other words, to ensure that the result you get in C0 is the same as the result you get with true integer arithmetic).

```
int safe_add(int a, int b)
/*@requires
```

```
@*/
{ return a + b; }
```

What about multiplication? For the sake of simplicity, you can assume both numbers are non-negative.

```
int safe_mult(int a, int b)
/*@requires a >= 0 && b >= 0 &&
```


## @*/

\{ return a * b; \}

## ARGB representation of color

We usually use 32-bit integers in C0 to represent a single integer. However, it's possible to use the bits in other ways: as 32 separate Boolean values or as 4 separate 8 -bit numbers in the range $[0,255)$. This lets us represent a color (red, green, and blue intensities, plus transparency or "alpha"), as 32-bit C0 integer.

| Sample Length: |  |  | 8 |  |  |  |  |  |  | 8 | 8 |  |  |  |  |  |  | 8 |  |  |  |  |  |  | 8 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Channel Membership: |  |  | Ph |  |  |  |  |  |  | Re | d |  |  |  |  |  |  | ree |  |  |  |  |  |  | u |  |  |  |
|  |  |  |  | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Bit Number: | 30 | 29 | 27 | 726 | 25 | 24 | 23 | 22 | 21 | 20 | 19 | 181 | 171 | 161 | 151 | 141 | 131 | 211 | 10 | 9 | 8 | 7 | 6 | 5 |  |  |  | 0 |

