15-122: Principles of Imperative Computation, Spring 2024 Written Homework 11

Due on Gradescope: Monday 1st April, 2024 by 9pm

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Section:		

This written homework provides practice with C features such as arrays, undefined behaviors, and stack allocation.

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- Kami, Adobe Acrobat Online, or *DocHub*, some web-based PDF editors that work from anywhere.
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Question:	1	2	3	4	Total
Points:	3.5	2.5	4	5	15
Score:					

1. Accessing Data on the Stack

In this question, we assume the usual 2's complement implementation of unsigned and signed **char** (8 bits, one byte), **short** (16 bits, two bytes), and **int** (32 bits, four bytes). Pointers are 64 bits, eight bytes. We also assume that numbers are represented in memory with their most significant bits at lower addresses.

During execution, the local variables of a function are stored on the **system stack**. In particular, so are stack-allocated arrays and structs. Each item on the stack starts at a unique address, but each address contains exactly one byte of data. Therefore the value of a variable will often span multiple bytes.

Assume the function being executed contains the following lo-Var. Addr. Stack cal declarations:

```
// a stack-allocated array
char cs[] = {'a', 'b', 'c', 'd'};
// an int
int i = 0x01ABCDEF;
```

The compiler decides on the layout of their component bytes onto the stack. Assume it arranged them as shown on the right of this text: the variable cs starts at address 0xC4 which contains 0x61 (the ASCII value of 'a'), the variable i starts at address 0xC8 where it stores the 8 most significant bits of its value.

Notice that while cs and i look similar on the stack, cs is an array — its values are read by indexing — whereas i is an integer with value 0x01ABCDEF.

. . . 0xCB EF CD 0xCA 0xC9 AB 0xC8 01 64 0xC7 0xC6 63 62 0xC5 61 CS 0xC4

i.

1.1 For the following program, draw the contents of the stack after the code shown is executed. Assume the compiler arranges the bytes as in the previous example. Leave any byte that does not contain program data blank.

char A[5];	Var. Addr. Stack
A[0] = '1'; A[1] = '5';	
A[2] = '1'; A[3] = '2';	0×BE
A[4] = '2';	0×BD 0×BC
The ASCII value of the character '1' is 0x31.	0×BB
What is the value of &A[1]?	0×BA
Fill in the correct type so that *x == A[3]: x = &A[3];	0×B9 A 0×B8

1pt

1.2 Assume we are given the function

bool parse(char *x, int *p); // Returns true iff parse succeeds

which reads the string representation x of a decimal number (e.g., "15122") and stores the corresponding integer (here 15122) at p. If x is not a decimal number (e.g., "hello!"), it returns false — it returns true to signal a successful parse.

Now we want to implement the function test_parse which calls parse on its input and prints whether the parse was successful or not. Before we knew about getting the address of a variable, we would have implemented test_parse as follows:

```
void test_parse(char *x) {
  REQUIRES(x != NULL);
  int *i = xmalloc(sizeof(int));
  if (parse(x, i)) printf("Success: %d!\n", *i);
  else printf("Failure.\n");
  free(i);
}
```

In particular, we had to allocate space for the variable i so that parse had somewhere to write the result of a successful parse.

But now that we know how to get the address of a variable, we can reimplement test_parse so that it does not heap-allocate i (and risk leaking memory). Complete the following code:

```
void test_parse_lean(char *x) {
  REQUIRES(x != NULL);
  int i;
}
```

1.3 Assume the string "410", which resides at address 0xA7,

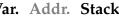
is passed into test_parse_lean so that, just before the call to parse, the contents of the stack are displayed to the

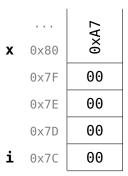
right of this text. For succinctness, we wrote the memory

address of the string "410" as 0xA7 on the stack drawing

— an actual address would span 8-bytes of stack space.

Var. Addr. Stack





What are the hex values of the arguments passed to parse? <i>Hint: recall what the values of pointer types are.</i>	Var. Addr. St	ack
parse(,)	 X 0×80	0×A7
What are the contents of the stack when parse has	0×7F	
finished executing? Write all values in hex in the dia- gram to the right of this text.	0×7E 0×7D	
<u> </u>	i 0x7C	

- 0.5pts
- **1.4** One of the dangers of using the address-of operator is that stack addresses get reused. Consider a version of test_parse_lean that returns the address of the parsed value.

```
int *test_parse_weird(char *x) {
    REQUIRES(x != NULL);
    int i;
    // Your code from test_parse_lean goes here
    return &i;
}
int main() {
    int *n1 = test_parse_weird("122");
    int *n2 = test_parse_weird("150");
    ASSERT(*n1 != *n2); // we expect n1 and n2 to differ
    return 0;
}
```

Since we don't know anything about where i appears on the stack, both calls to test_parse_weird could end up putting i at the same address (i.e., the values would be overwritten). Answer the following questions under this assumption.

If n1 == 0xB0, the value of n2 is ______ The argument of ASSERT evaluates to ______ (true or false?)

2. Undefined Behavior

C is difficult to program in because it is so permissive. Many things which are errors in C0 are undefined in C, so the behavior can change even between runs.

Consider the following program, where we assume **int**s are 4 bytes long and pointers are 8 bytes long:

```
1 #include <stdio.h>
2 #include <stdlib.h>
3
4 int main() {
    int *n = malloc(2 * sizeof(int));
5
    *n = 15122;
6
7
    int *A;
8
   *A = 6;
9
10
    if (*n == 15122) printf("Yay!\n");
11
    else printf("Blasphemy\n");
12
13
    free(n);
14
    return 0;
15
16 }
```

Assume that when the execution reaches line 7, the heap looks as follows:

			n	•]														
OxAD	OxA1	OxA2	0xA3	OxA4	OxA5	0xA6	0xA7	0xA8	OxA9	OxAA	OxAB	OxAC	0xAD	Oxae	0xAF	0xB0	OxB1	OxB2	OxB3

The grayed out region is allocated by you on line 5 — i.e., malloc returned 0xA4. Access to the red region is forbidden (doing so will result in a segmentation fault). Although executing line 9 is undefined behavior, you may assume it succeeds unless it writes to the red region.

2.1 Since local variables are not automatically initialized, A could be anything. For each of the following possible values of A, state whether the program runs normally, prints "blasphemy", or throws an error. Mark only one outcome for each row.

A == 0×A0:	O prints "Yay!"	◯ prints "Blasphemy"	Segfaults
A == 0xA4:	O prints "Yay!"	◯ prints "Blasphemy"	◯ segfaults
A == 0xA6:	O prints "Yay!"	◯ prints "Blasphemy"	◯ segfaults
A == 0xA8:	O prints "Yay!"	◯ prints "Blasphemy"	◯ segfaults
A == 0xAC:	O prints "Yay!"	◯ prints "Blasphemy"	◯ segfaults
A == 0×AE:	O prints "Yay!"	◯ prints "Blasphemy"	◯ segfaults

For the next two task, you will need to compile this code and run it with Valgrind. You will need to use the following compilation command:

gcc -g file.c

(without additional flags) where file.c is the name of the file where you saved this code. Read the guide to success on **How to use Valgrind** on o interpret the output of Valgrind.

0.5pts

0.5pts

1.5pts

2.2 Valgrind will report two error messages with this code even if the program runs *correctly*. What is the problem in the code that causes both error messages?

Problem:			

2.3 Using this error message, fix the error in file.c. Then, change line 14 to free(n+1). What is the *exact error message* that Valgrind displays now? Explain your answer.

Valgrind error message: _____

because

Always use Valgrind!

3. Strings in C

Unlike C0, C does not have a **string** type. Rather, strings in C are arrays of **char**'s that end in the NUL character, the character with ASCII value 0, which we write as $' \setminus 0'$. They can live in three places:

- 1. In the **DATA** segment of memory. Space is automatically created for them, it is undefined behavior to write to them (they are "read-only"), and they are automatically NUL-terminated. You create them when you write string literals, such as "functions are pointers", in your code.
- 2. On the **heap**. Space needs to be provided for them using malloc or similar, and it needs to be freed.
- 3. On the **stack**, as stack allocated arrays. Space for them is created when entering a function and disappears when returning from it.

As you might anticipate, strings can lead to a number of errors! The tasks in this question will explore the ways things can go wrong.

3.1 The following function writes A to res, followed by these same characters in reverse. For example, if A is "honk!", then res will be "honk!!knoh".

```
void mirror(char *A, char *res) {
    REQUIRES(A != NULL);
2
    REQUIRES(res != NULL);
3
4
    int n = strlen(A);
5
    for (int i = 0; i < n; i++) {</pre>
6
      res[i] = A[i];
7
      res[n+i] = A[n-i-1];
8
    }
9
    res[2*n] = ' \setminus 0';
10
11 }
```

Depending on what string is passed to mirror as A and what we do with our memory, we could run into a number of common issues:

- We could read a value out of bounds or read from memory we have already freed. These are *invalid reads*.
- We could write to a value out of bounds, write to memory we have already freed, or write to read-only memory. These are *invalid writes*.
- We could free a value we have already freed or free a pointer we did not allocate. These are *invalid frees*.

1.5pts

Here are some test cases using mirror. Determine whether they could result in an invalid read, an invalid write, or an invalid free. For each one, check all boxes that apply. If none do, check "ALL GOOD". Remember to account for the NUL-terminator!

When determining the effects of each statement, assume the previous ones (except the setup code before the first call to mirror) did not run and undefined behavior does not propagate to later function calls on the same line.

The library function strncpy(s1,s2,n) copies the first n characters of s2 into s1.

```
int main() {
```

```
char A[] = "hi!";
2
    char B[] = {'h', 'i','!'};
   char *C = "hi!":
    char *D = xmalloc(sizeof(char) * 4);
   strncpy(D, C, 4);
   char *E = D;
   char *a = xmalloc(sizeof(char) * 7);
    char b[7];
10
   char c[] = "hi";
11
    char *d = "goodbye";
12
    char *e = xmalloc(sizeof(char) * 12);
13
14
    11
                                       Invalid Invalid Invalid
                                                                      ALL
15
                                                  write
                                                                     GOOD
    11
                                        read
                                                            free
16
   mirror(A, a);
17
   mirror(B, b);
18
   mirror(C, c);
19
   mirror(D, d);
20
   mirror(E, e);
21
   mirror(A, b); free(b);
22
   mirror(C, e); free(e);
23
   free(d); mirror(E, c);
24
   free(a); mirror(B, d);
25
   free(E); mirror(D, a); free(D);
26
    return 0;
27
28 }
```

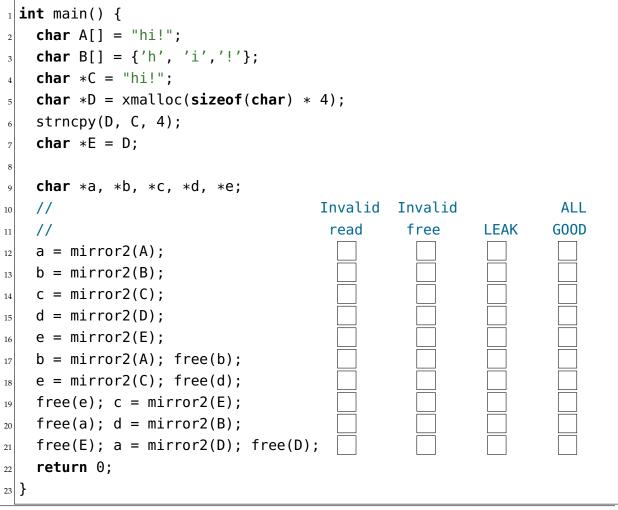
3.2 We can't fix all of these problems, but we can eliminate some by returning the result. Complete the code of mirror2 so that it cannot perform invalid writes.

```
char *mirror2(char *A) {
    REQUIRES(A != NULL);
    int n = strlen(A);
    char *res = xcalloc(sizeof(char), _____);
    return _____;
}
```

1.5pts

1pt

3.3 This version prevents invalid writes, but it introduces more possible leaks. For each line, mark it if it allocates memory that gets leaked when the program reaches line 22. Assume that all earlier lines ran but ignore any undefined behavior on them.



4. Casting C Things

Since C has integers of different sizes, we often find it necessary to cast between them. In this question, we will make the following assumptions:

- whether signed or unsigned, a **char** has 8 bits (one byte), a **short** has 16 bits (two bytes), and an **int** has 32 bits (four bytes);
- signed numbers are implemented using two's complement (or equivalently the code was compiled with the fwrapv flag).
- 4.1 Given the declarations

int sign_int = -7; unsigned int un_int = 40125; signed char sign_char = -114; unsigned char un_char = 221;

fill in the table below. In this task, fully simplify decimal values and always write the leading zeros in the hex representation, so that an **int** is 8 hex digits, a **short** 4, and a **char** 2. If the result of a cast is implementation-defined, assume the common rules seen in class but flag this issue by writing an X in column **ID**.

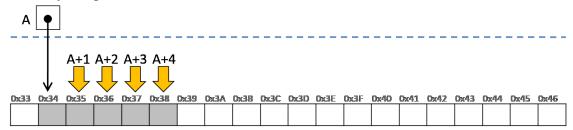
	C expression	Decimal value	Hexadecimal	ID
1.	sign_int	- 7	0xFFFFFF9	
2.	(unsigned int)sign_int			
3.	un_int	40125		
4.	(int)un_int			
5.	(unsigned short)un_int			
6.	<pre>(short)(unsigned short)un_int</pre>			
7.	sign_char	-114		
8.	(int)sign_char			
9.	(unsigned int)(int)sign_char			
10.	un_char	221		
11.	(signed char)un_char			
12.	(int)(signed char)un_char			
13.	(unsigned int)(int)(signed char)un_char			
14.	(unsigned int)un_char			

2pts

1pt

4.2 An array's address is its starting byte. The type of an array tells the compiler how many bytes to read when accessing an element.

Consider the array A represented in memory by the shaded cell in the following memory diagram:



What are the possible types of the elements of A? Mark all that apply.

char	unsigned char	short	unsigned short
int	unsigned int	long	unsigned long
int8_t	int16_t	int32_t	int64_t
uint8_t	uint16_t	uint32_t	uint64_t

Ananya wants to convert an array A from the previous task so that each element has type **int**. For example, if A contains the characters {0x01 0xF1 0x80 0x02 0x36}, the new array should contain {0x0000001 0x000000F1 0x00000080 0x00000002 0x00000036} — it's the same as A but each element is now an **int** instead of a **char**.

Ananya's friend, Grace, suggests that Ananya converts her array by defining the following:

int *B = (int*)A;

Clearly indicate B, B+1, B+2, B+3 and B+4 in the following memory diagram.

0x33 0x34 (0x35 0x3	36 0x37	0x38	0x39	0x3A	0x3B	0x3C	0x3D	0x3E	0x3F	0x40	0x41	0x42	Ox43	0x44	0x45	0x46

1pt

1pt

- **4.3** Grace's approach doesn't quite work. Help Ananya out by writing a function that converts her array correctly. Fill in the missing types appropriately and make all casts explicit. As you do so, recall C's casting rules for numerical types:
 - If the new type can represent the value, the value is preserved.
 - If the new type can't represent the value,
 - if the new type is unsigned,
 - * if the new type is smaller or the same, the least significant bits are retained.
 - * if the new type is bigger, the bits are sign-extended.
 - if the new type is signed, the result is implementation-defined.

You may not use bitwise operators.

As earlier, if A contains the characters $\{0x01 \ 0xF1 \ 0x80 \ 0x02 \ 0x36\}$, the new array should contain $\{0x00000001 \ 0x000000F1 \ 0x00000080 \ 0x00000002 \ 0x00000036\}$.

```
int *char_array_to_ints(signed char *A, size_t num_elems) {
}
```

4.4 We want to convert an array of numbers of unspecified type (they could be char's, uint64_t's, ...) to an array of int's. How to go about it? As long as we know the size of the array elements and how to turn this many bytes to an int, we can perform this conversion. Complete the code of the following function that does precisely this! You will want to use the fact that sizeof(char) == 1. Make all casts explicit.

}