Database Applications (15-415)

SQL-Part III &
Storing Data: Disks and Files- Part I
Lecture 8, February 5, 2014

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Today...

- Last Session:
  - Standard Query Language (SQL)- Part II

- Today’s Session:
  - Standard Query Language (SQL)- Part III
  - Storing Data: Disks and Files- Part I

- Announcements:
  - Project 1 is now posted. Due on Feb 18 by midnight
  - Quiz I is on Monday 10, 2014 (all topics included except today’s data storage topic)
  - PS2 is due on Feb 07 by midnight
Outline

NULL values and Join Variants

Complex Integrity Constraints and Triggers

Java Database Connectivity

Storing Data: Disks and Files (Briefly)
NULL Values

- Column values can be *unknown* (e.g., a sailor may not yet have a rating assigned)

- Column values may be *inapplicable* (e.g., a maiden-name column for men!)

- **NULL** values can be used in such situations

- However, NULL values complicate many issues!
  - Comparing NULL to a valid value returns unknown
  - Comparing NULL to a NULL returns unknown
NULL Values

- Considering a row with rating = NULL and age = 20; How does it compare with the following Boolean expressions?
  - Rating = 8 OR age < 40 ➔ TRUE
  - Rating = 8 AND age < 40 ➔ unknown

- In general, what about?
  - NOT unknown ➔ unknown
  - True OR unknown ➔ True
  - False OR unknown ➔ unknown
  - False AND unknown ➔ False
  - True AND unknown ➔ unknown
NULL Values

- Considering a row with rating = NULL and age = 20; How does it compare with the following Boolean expressions?
  - Rating = 8 OR age < 40 ➔ TRUE
  - Rating = 8 AND age < 40 ➔ unknown

- In general, what about?
  - NOT unknown ➔ unknown
  - True OR unknown ➔ True
  - False OR unknown ➔ unknown
  - False AND unknown ➔ False
  - True AND unknown ➔ unknown

*Three-Valued Logic!*
Inner Joins

- Tuples of a relation that do not match some row in another relation (according to a join condition $c$) do not appear in the result
  - Such a join is referred to as “Inner Join” (so far, all inner joins)

```sql
select ssn, c-name
from takes, class
where takes.c-id = class.c-id
```

Equivalently:

```sql
select ssn, c-name
from takes join class on takes.c-id = class.c-id
```
An Example of Inner Joins

- Find all SSN(s) taking course s.e.

<table>
<thead>
<tr>
<th>TAKES</th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>SSN</td>
<td>c-id</td>
<td>grade</td>
</tr>
<tr>
<td>123</td>
<td>15-413</td>
<td>A</td>
</tr>
<tr>
<td>234</td>
<td>15-413</td>
<td>B</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>CLASS</th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>c-id</td>
<td>c-name</td>
<td>units</td>
<td></td>
</tr>
<tr>
<td>15-413</td>
<td>s.e.</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>15-412</td>
<td>o.s.</td>
<td>2</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>SSN</th>
<th>c-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>123</td>
<td>s.e</td>
</tr>
<tr>
<td>234</td>
<td>s.e</td>
</tr>
</tbody>
</table>

**o.s.: gone!**
Outer Joins

- Tuples of a relation that do not match some row in another relation (according to a join condition $c$) can still appear exactly once in the result
  - Such a join is referred to as “Outer Join”
  - Result columns will be assigned NULL values

```sql
select ssn, c-name
from takes outer join class
on takes.c-id=class.c-id
```
An Example of Outer Joins

- Find all SSN(s) taking course s.e.

<table>
<thead>
<tr>
<th>SSN</th>
<th>c-id</th>
<th>grade</th>
</tr>
</thead>
<tbody>
<tr>
<td>123</td>
<td>15-413</td>
<td>A</td>
</tr>
<tr>
<td>234</td>
<td>15-413</td>
<td>B</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>SSN</th>
<th>c-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>123</td>
<td>s.e.</td>
</tr>
<tr>
<td>234</td>
<td>s.e.</td>
</tr>
<tr>
<td>null</td>
<td>o.s.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>c-id</th>
<th>c-name</th>
<th>units</th>
</tr>
</thead>
<tbody>
<tr>
<td>15-413</td>
<td>s.e.</td>
<td>2</td>
</tr>
<tr>
<td>15-412</td>
<td>o.s.</td>
<td>2</td>
</tr>
</tbody>
</table>
Joins

- The general SQL syntax:

```sql
select [column list]
from table_name
[inner | {left | right | full} outer] join table_name
on qualification_list
```

<table>
<thead>
<tr>
<th>Outer Join Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Left Outer Join</td>
<td>A rows without a matching B row appear in the result</td>
</tr>
<tr>
<td>Right Outer Join</td>
<td>B rows without a matching A row appear in the result</td>
</tr>
<tr>
<td>Full Outer Join</td>
<td>Both A and B rows without a match appear in the result</td>
</tr>
</tbody>
</table>
Outline

- NULL values and Join Variants
- Complex Integrity Constraints and Triggers
- Java Database Connectivity
- Storing Data: Disks and Files (Briefly)
Integrity Constraints - A Review

- An Integrity Constraint (IC) describes conditions that every legal instance of a relation must satisfy.

- Inserts/deletes/updates that violate IC’s are disallowed.

- ICs can be used to:
  - Ensure application semantics (e.g., sid is a key)
  - Prevent inconsistencies (e.g., sname has to be a string, age must be < 20)
Types of Integrity Constraints- A Review

- IC types:
  - Domain constraints
  - Primary key constraints
  - Foreign key constraints
  - General constraints
    - Useful when more general ICs than keys are involved
    - Can be specified over a single table and across tables
General Constraints Over a Single Table

- Complex constraints over a single table can be defined using `CHECK conditional-expression`

```sql
CREATE TABLE Sailors (sid INTEGER, sname CHAR (10), rating INTEGER, age REAL, PRIMARY KEY (sid), CHECK (rating >= 1 AND rating <= 10))
```

- A primary key constraint
- A general constraint
- A domain constraint
How can we enforce that “Interlake” boats cannot be reserved?

```
CREATE TABLE Reserves (sid INTEGER, bid INTEGER, day DATE,
FOREIGN KEY (sid) REFERENCES Sailors,
FOREIGN KEY (bid) REFERENCES Boats,
CONSTRAINT noInterlakeRes,
CHECK ('Interlake' NOT IN (SELECT B.bname
FROM Boats B
WHERE B.bid = Reserves.bid)))
```
General Constraints Across Tables

- How can we enforce that the number of boats plus the number of sailors should not exceed 100?

CREATE TABLE Sailors (sid INTEGER, 
sname CHAR (10), 
rating INTEGER, 
age REAL, 
PRIMARY KEY (sid), 
CHECK (rating >= 1 AND rating <= 10) 
CHECK ( ((SELECT COUNT (S.sid) FROM Sailors S) + 
(SELECT COUNT (B.bid) FROM Boats B)) < 100))

What if the Sailors table is empty and we insert more than 100 rows into Boats?
How can we enforce that the number of boats plus the number of sailors should not exceed 100?

CREATE ASSERTION smallClub
CHECK
( (SELECT COUNT (S.sid) FROM Sailors S) + (SELECT COUNT (B.bid) FROM Boats B) < 100 )

ASSERTION is the right solution; not associated with either table!
New Domains

- Users can define new domains using the `CREATE DOMAIN` statement.

```sql
CREATE DOMAIN ratingval1 INTEGER DEFAULT 1
    CHECK ( VALUE >= 1 AND VALUE <= 10 )

CREATE DOMAIN ratingval2 INTEGER DEFAULT 1
    CHECK ( VALUE >= 1 AND VALUE <= 20 )
```

ratingval1 and ratingval2 CAN be compared!

Domain constraints will be always enforced (also for new domains)!
Distinct Types

- Users can define new distinct types using the CREATE TYPE statement

```sql
CREATE TYPE ratingtype1 AS INTEGER
```

```sql
CREATE TYPE ratingtype2 AS INTEGER
```

- ratingtype1 and ratingtype2 CANNOT be compared!

- Domain constraints will be always enforced (also for new types)!
Triggers

- A trigger is a *procedural* code that is automatically executed in response to certain *events* on a particular table or view in a database.

- Triggers can be activated either *before* or *after*:
  - Insertions
  - Deletions
  - Updates
A Trigger Example

- Set a timestamp field whenever a row in the takes table is updated

<table>
<thead>
<tr>
<th>TAKES</th>
</tr>
</thead>
<tbody>
<tr>
<td>SSN</td>
</tr>
<tr>
<td>-------</td>
</tr>
<tr>
<td>123</td>
</tr>
<tr>
<td>234</td>
</tr>
</tbody>
</table>

- **First**: we need to add our timestamp field

```
ALTER TABLE takes
ADD COLUMN updated TIMESTAMP
```
A Trigger Example

- Set a timestamp field whenever a row in the takes table is updated

<table>
<thead>
<tr>
<th>SSN</th>
<th>c-id</th>
<th>grade</th>
</tr>
</thead>
<tbody>
<tr>
<td>123</td>
<td>15-413</td>
<td>A</td>
</tr>
<tr>
<td>234</td>
<td>15-413</td>
<td>B</td>
</tr>
</tbody>
</table>

- **Second**: we need to create a function that sets the “updated” column with the current timestamp

```sql
CREATE FUNCTION update_col()
BEGIN
    NEW.updated = NOW();
    RETURN NEW;
END;
```
A Trigger Example

- Set a timestamp field whenever a row in the takes table is updated

<table>
<thead>
<tr>
<th>TAKES</th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>SSN</td>
<td>c-id</td>
</tr>
<tr>
<td>123</td>
<td>15-413</td>
</tr>
<tr>
<td>234</td>
<td>15-413</td>
</tr>
</tbody>
</table>

- Third: we need to Invoke update_col() when a row in the takes table is updated

CREATE TRIGGER update_takes_modtime
AFTER UPDATE ON takes
FOR EACH ROW
EXECUTE PROCEDURE update_col();

A row-level trigger; otherwise, it will be a statement-level trigger
Outline

- NULL values and Join Variants
- Complex Integrity Constraints and Triggers
- Java Database Connectivity
- Storing Data: Disks and Files (Briefly)
Java Database Connectivity

- SQL commands can be *embedded* in host language programs

- A popular data access technology which provides an API for querying and manipulating data in (any) storage system is called *Java Database Connectivity (JDBC)*

- Direct interactions with a DBMS occurs through a DBMS-specific *driver*

- A driver is a software program that translates JDBC calls into DBMS-specific calls
  - Drivers do not necessarily interact with a DBMS that understands SQL
  - Thus, a DBMS in JDBC’s parlance is usually referred to as *data source*
Establishing a Connection

- With JDBC, a database is represented by a URL

- With PostgreSQL™, this takes one of the following forms:
  - `jdbc:postgresql:database`
  - `jdbc:postgresql://host/database`
  - `jdbc:postgresql://host:port/database`

- To connect to a database, a Connection instance from JDBC can be used

  ```java
  Connection db = DriverManager.getConnection(url, username, password);
  ```
Establishing a Connection

- A number of additional properties can be used to specify additional driver behavior specific to PostgreSQL™

```java
String url = "jdbc:postgresql://localhost/test";
Properties props = new Properties();
props.setProperty("user","Hammoud");
props.setProperty("password","secret");
props.setProperty("ssl","true");
Connection conn = DriverManager.getConnection(url, props);
```

Equivalently:

```java
String url = "jdbc:postgresql://localhost/test?user=Hammoud&password=secret&ssl=true";
Connection conn = DriverManager.getConnection(url);
```
Establishing a Connection

- Putting it all together, you can create the following function:

```java
public Connection getConnection() throws SQLException {
    String url = "jdbc:postgresql://localhost/test";
    Properties props = new Properties();
    props.setProperty("user","Hammoud");
    props.setProperty("password","secret");
    props.setProperty("ssl","true");
    Connection conn = DriverManager.getConnection(url, props);
    System.out.println("Connected to database");
    return conn;
}
```
Creating Tables

- Assume the following students table:

<table>
<thead>
<tr>
<th>Sid</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Hammoud</td>
</tr>
<tr>
<td>2</td>
<td>Esam</td>
</tr>
</tbody>
</table>

**SQL:**

```sql
CREATE TABLE students(
    sid INTEGER,
    name CHAR(30),
    PRIMARY KEY (sid)
);
```

**JDBC:**

```java
public void createTable() throws SQLException {
    String createT = "create table students (sid INTEGER, " +
                     "name CHAR(30) " +
                     "PRIMARY KEY (sid))";
    Statement stmt = null;
    try {
        stmt = conn.createStatement();
        stmt.executeUpdate(createT);
    } catch (SQLException e) {
        e.printStackTrace(e);
    }
    finally { if (stmt != null) { stmt.close(); } }
}
```
Populating Tables

- Assume the following students table:

<table>
<thead>
<tr>
<th>Sid</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Hammoud</td>
</tr>
<tr>
<td>2</td>
<td>Esam</td>
</tr>
</tbody>
</table>

**SQL:**
- `INSERT INTO students values (1, 'Hammoud')`
- `INSERT INTO students values (2, 'Esam')`

**JDBC:**
```java
public void populateTable() throws SQLException {
    Statement stmt = null;
    try {
        stmt = conn.createStatement();
        stmt.executeUpdate("insert into students values(1, 'Hammoud')");
        stmt.executeUpdate("insert into students values(2, 'Esam')");
    } catch (SQLException e) {} 
    finally { if (stmt != null) { stmt.close(); } }
}
```
Querying Tables

- Assume the following students table:

<table>
<thead>
<tr>
<th>Sid</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Hammoud</td>
</tr>
<tr>
<td>2</td>
<td>Esam</td>
</tr>
</tbody>
</table>

**SQL:** SELECT sid, name from students

**JDBC:**

```java
public static void viewTable() throws SQLException {
    Statement stmt = null;
    String query = "select sid, name from students";
    try {
        stmt = conn.createStatement();
        ResultSet rs = stmt.executeQuery(query);
        while (rs.next()) {
            int sID = rs.getInt("sid");
            String sName = rs.getString("name");
            System.out.println(sName + "\t" + sID);
        }
    } catch (SQLException e) {} finally { if (stmt != null) { stmt.close(); } }
}
```

A “cursor” the points to one row of data at a time

Columns retrieved by names
Querying Tables

Assume the following students table:

<table>
<thead>
<tr>
<th>Sid</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Hammoud</td>
</tr>
<tr>
<td>2</td>
<td>Esam</td>
</tr>
</tbody>
</table>

**SQL:** SELECT sid, name from students

**JDBC:**

```java
public static void viewTable() throws SQLException {
    Statement stmt = null;
    String query = "select sid, name from students";
    try {
        stmt = conn.createStatement();
        ResultSet rs = stmt.executeQuery(query);
        while (rs.next()) {
            int sID = rs.getInt(1);
            String sName = rs.getString(2);
            System.out.println(sName + "\t" + sID);
        }
    } catch (SQLException e) {} finally { if (stmt != null) { stmt.close(); } }
}
```

**OR:** Columns retrieved by numbers
Cursor Methods

- Methods available to move the cursor of a result set:
  - `next()`
  - `previous()`
  - `first()`
  - `Last()`
  - `beforeFirst()`
  - `afterLast()`
  - `relative(int rows)`
  - `absolute(int row)`

By default, you can call only `next()`! More on this shortly...
Updating Tables

- By default, ResultSet objects cannot be updated, and their cursors can only be moved forward

- ResultSet objects can be though defined to be *scrollable* (the cursor can move backwards or move to an absolute position) and *updatable*

```java
public void modifyStudents() throws SQLException {
    Statement stmt = null;
    try {
        // stmt = con.createStatement(); */
        stmt = con.createStatement(ResultSet.TYPE_SCROLL_SENSITIVE,
                                   ResultSet.CONCUR_UPDATABLE);
        ResultSet uprs = stmt.executeQuery("SELECT * FROM students");
        while (uprs.next()) {
            String old_n = uprs.getString("name");
            uprs.updateString("name", "Mohammad" + old_n);
            uprs.updateRow();
        }
    } catch (SQLException e) {} finally { if (stmt != null) { stmt.close(); } }
}
```
Result Set Types

- **TYPE_FORWARD_ONLY** (the default)
  - The result set is **not** scrollable

- **TYPE_SCROLL_INSENSITIVE**
  - The result set is scrollable
  - The result set is **insensitive** to changes made to the underlying data source while it is open

- **TYPE_SCROLL_SENSITIVE**
  - The result set is scrollable
  - The result set is **sensitive** to changes made to the underlying data source while it is open
Result Set Concurrency

- The concurrency of a ResultSet object determines what level of update functionality is supported

- Concurrency levels:
  - `CONCUR_READ_ONLY` (the default)
    - The result set **cannot** be updated
  - `CONCUR_UPDATABLE`
    - The result set **can** be updated
Prepared Statements

- JDBC allows using a PreparedStatement object for sending SQL statements to a database

- This way, the same statement can be used with different values many times

```java
String sql = "INSERT into students values (?, ?)";
PreparedStatement ps = conn.prepareStatement(sql);
ps.clearParameters();
ps.setInt(1, 111);
ps.setString(2, "Hammoud");
int numRows1 = ps.executeUpdate();

ps.setInt(1, 222);
ps.setString(2, "Esam");
int numRows2 = ps.executeUpdate();
```

More about JDBC in the recitations!
Outline

- NULL values and Join Variants
- Complex Integrity Constraints and Triggers
- Java Database Connectivity
- Storing Data: Disks and Files (*Briefly*)
DBMS Layers

 Queries

 Query Optimization and Execution

 Relational Operators

 Files and Access Methods

 Buffer Management

 Disk Space Management

 Transaction Manager

 Lock Manager

 Recovery Manager

 DB

 Today and Next Lectures
The Memory Hierarchy

- Storage devices play an important role in database systems

- How systems arrange storage?

More expensive, but faster!

Less expensive, but slower!

Disk

Storage devices play an important role in database systems

- More expensive, but faster!
- Less expensive, but slower!
Where to Store Data?

- Where does DBMS store information?
  - DBMSs store large amount of data (e.g., Big Data!)

- Buying enough memory to store all data is prohibitively expensive (let alone that memories are volatile)

- Thus, databases are usually stored on disks (or tapes for backups)
But, What Will Do With Memory?

- Data must be brought to memory to be processed!
  - **READ:** transfer data from disk to main memory (RAM)
  - **WRITE:** transfer data from RAM to disk

- I/O time dominates the time taken for database operations!
- To minimize I/O time, it is necessary to store and locate data strategically
Magnetic Disks

- Data is stored in disk blocks
- Blocks are arranged in concentric rings called tracks
- The arm assembly is moved in or out to position a head on a desired track
- Each track is divided into arcs called sectors (whose size is fixed)
- The block size is a multiple of sector size
- The set of all tracks with the same diameter is called cylinder
Accessing a Disk Block

- What is I/O time?
  - The time to move the disk heads to the track on which a desired block is located
  - The waiting time for the desired block to rotate under the disk head
  - The time to actually read or write the data in the block once the head is positioned
Accessing a Disk Block

- What is I/O time?
  - Seek Time
  - Rotational Time
  - Transfer Time

- I/O time = seek time + rotational time + transfer time
Implications on DBMSs

- Seek time and rotational delay dominate!

- Key to lower I/O cost: reduce seek/rotation delays!

- How to minimize seek and rotational delays?
  - Blocks on same track, followed by
  - Blocks on same cylinder, followed by
  - Blocks on adjacent cylinder
  - Hence, _sequential_ arrangement of blocks in a file is a big win!

Who can take care of that?
Storing Data: Concluding Remarks

- Disks provide cheap, non-volatile storage for DBMSs
- However, data must be in memory for the DBMS to operate on it
- I/O time dominates!
- The cost depends on the locations of pages on disk
- It is important to arrange data sequentially to minimize *seek* and *rotation* delays
Next Class

Quiz I &

Storing Data: Disks and Files

(Cont’d)