

15-440
Distributed Systems
Recitation 7

Tamim Jabban

Project 2

- Involves *building on your Project 1 Distributed File System (DFS): FileStack*
- *P2_StarterCode*: Copy files into your P1 folder
- *Due date*: November 1st

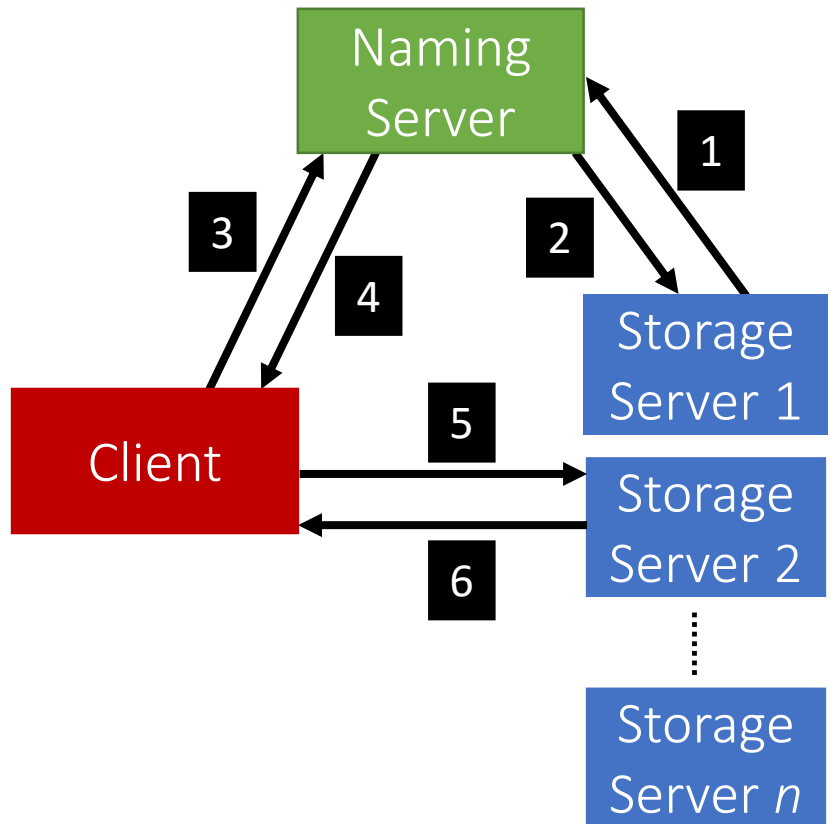
Project 1: Recap

- Applied the knowledge of client-server communication and **Remote Method Invocation (RMI)** to build a Distributed File System denoted as FileStack
- Employed **stubs and skeletons** to mask communication, thereby transparently locating and manipulating files stored remotely at a cluster of machines

Entities & Architecture

- Storage Servers (SSs)
 - Each SS stores physically files to share in a directory (denoted as temporary directory) in its local file system
- Naming Server (NS)
 - Stores metadata about all shared files in the form of a mapping from filenames to storage servers (like DNS)
- Clients
 - Perform operations on files (e.g., write, read etc.)
- Architecture
 - Based on client-server architecture

Communication b/w Entities



Request-Reply
Communication Paradigm

- 1 Registration
- 2 Duplicate Files, Create, Delete
- 3 CreateFile, CreateDirectory, IsDirectory, Delete, List, GetStorage
- 4 Results, Storage Stub
- 5 Read, Write, Size
- 6 Results (*of Read, Write, Size*)

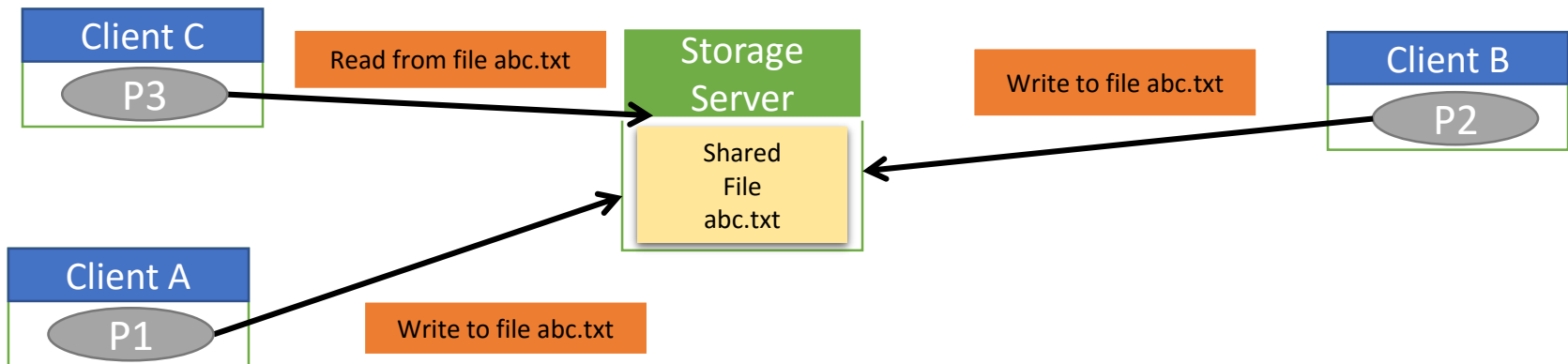
File Correctness & Consistency

- Did we allow multiple clients to write on a file?

Yes!

- Did we allow a client to read a file under modification?

Yes!



Project 2 Objectives

1. Devise and apply a **synchronization algorithm** that:
 - achieves *correctness* while sharing files
 - and ensures *fairness* to clients.
2. Devise and apply a **replication algorithm** that:
 - achieves load-balancing among storage servers
 - and ensures consistency of replicated files.

Project 2 Objectives

1. Logical Synchronization of Readers and Writers
2. Devise and apply a replication algorithm that:
 - achieves load-balancing among storage servers
 - and ensures consistency of replicated files.

Mutual Exclusion

1. Reader:

- Reader is a Client who wishes to read a file at a SS
- Reader first requests a [read/non-exclusive/shared lock](#)

2. Writer:

- Writer is a Client who wishes to write to a file at a SS
- Writer first requests a [write/exclusive lock](#)

3. Order:

- Readers and writers are queued and served in the [FIFO](#) order

Read Locks

- Readers do not modify contents of a file/directory
- Readers request the NS for read locks before reading files
- Readers unlock files once done
- Multiple readers can acquire a read lock simultaneously

Write Locks

- Writers can modify contents of files/directories
- Writers request the NS for write locks before reading/writing to files
- Writers unlock files once done
- Only one writer can acquire a write lock at a time

Write Locks

- NS grants a write lock on a file if:
 - No reader is currently reading the file
 - No writer is currently writing to the file
- Assume a writer requests a write lock for `project2.txt`
`/FileStack/users/student1/work/project2.txt`
- NS applies read locks on all the directories in the path to prevent modifications
- NS then grants a write lock to the requestor of `project2.txt`

Service Interface

- Two new operations available to Clients:
 - LOCK(path, read/write)
 - UNLOCK(path, read/write)

Project 2 Objectives

1. Devise and apply a synchronization algorithm that:
 - achieves *correctness* while sharing files
 - and ensures *fairness* to clients.
2. Dynamic Replication of Files

Why Replicate?

- In our DFS, we'll have two kinds of Files:
 - Files that have a lot of requests
 - These are denoted as "*hot-files*"
 - Files that are very rarely accessed
 - These are denoted as "*cold-files*"
- To achieve load-balancing, we can replicate "*hot-files*" onto other SSs

How Many Replicas?

- To measure file “hot-ness,” the NS can keep track of the number of requests to a file:
 - *num_requesters*: number of read requests to a file
- *To scale replicas linearly with the increase of num_requests:*
 - $num_replicas = \alpha * num_requesters$

How Many Replicas?

- However, we need to limit the number of replicas:
 - $num_replicas = \min(\alpha * num_requesters, upper_bound)$
- This is still too sensitive/fine-grained:
 - $num_requests_coarse$: $num_requests$ rounded to the nearest multiple of 20
 - $num_replicas = \min(\alpha * num_requests_coarse, replica_upper_bound)$

When to Replicate?

- NS would want to store *num_requests* as file metadata
- However, how can we determine and in turn update *num_requests* over time?
 - We know that Clients invoke read operations on storage servers
 - Therefore, every “read” lock request from a client is deemed as a read operation
 - Afterward, NS increments *num_requests*
 - Reevaluate *num_replicas*

How can we Replicate?

- NS first elects SSs to store the replicas
- NS commands each elected SS to copy the file from the original SS
- Therefore, the metadata of a file now includes *a set of SSs* instead of a single SS

How to Update Replicas

- When a Client requests a write lock on a file:
 - It causes the NS to *invalidate* all the replicas except the locked one
- Invalidation is achieved by **commanding those SSs hosting replicas to delete the file**
- When the Client unlocks the file, the NS commands SSs to copy the modified file

The Command Interface

- One new operation available to the NS:
 - `Copy(path, StorageStub)`