15-440
Distributed Systems

Recitation 5

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& Ammar Karkour
Announcements

Grades for Design Report – Today

Problem Set 2
Due: Sep. 26th (Tuesday)
Outline

• RPC Failure Sources
• Handling Different Types or RPC Failures
  • Communication Failures
  • Client-Side Failures
  • Server-Side Failures
• Calculating RPC Time
• Exercises
• Feedback Reflection
• Project 1 & PS2 Questions
Remote Procedure Calls (RPC) Failures

- Failures should be taken into consideration when RPC calls are made.

**Where Can Failures Happen & How to Handle Them**
RPC Failures – Where Can Failure Happen?

- Request gets lost
- Request received
- Server crashes
- Request received
- Server crashes
- Work done
- Reply sent
- Client crashes before receiving result
- Orphan work

Source: https://cseweb.ucsd.edu/classes/sp16/cse291-e/applications/ln/lecture3.html
RPC Failures – Types

Can be Categorized Into Three Types:
- Communication Failures
- Client Failures
- Server Failures

- Request gets lost
- Request Received
- Server Crashes
- Request Received
- Work Done
- Server Crashes
- Request Received
- Work Done
- Reply gets Lost
- Client Crashes Before Receiving Result

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Handling RPC Failures - Communication

• If the client is unable to find the server,
  • The client should return an error message.

• If client and server both are running but something goes wrong on the network and....
  
  RPC request packet doesn’t reach the server for processing or
  the server sent the reply back to client but client never received the reply.

  Retransmission is Needed!

  **Who should handle Retransmission?**

  Depends on the Transport Layer Protocol that RPC is layered on top of
Handling RPC Failures - Communication

• Who should Handle Retransmission of Lost Requests/Responses?

• If the **RPC system is running over UDP**, then RPC should handle it using a timeout mechanism:
  
  • The client should implement a timeout mechanism. If the client doesn’t receive response from server before the timeout then client should resend the RPC request.
  
  • Also, server in some RPC semantics (as we will see) will set a timeout waiting for client’s ACK after it sends the response.

If the **RPC system is running over TCP**, then the error correction task should be offloaded to the **TCP layer**.

A packet that is sent will arrive at the other end eventually.

**But this doesn’t handle Failures at End Points!**
Handling RPC Failures - Communication

When is it better to use TCP & UDP for RPC?

**UDP: Faster but Unreliable**
- in LAN,
- For certain kinds of Apps; e.g.:
  - Live streaming Apps
  - Apps that require fast communication (VoIP, online games, etc.)

**TCP: Slower but Reliable**
- in WAN,
- Apps that require full transmission of data (e.g. online banking)
- Large messages (as TCP can fragment)
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RPC Failure – Client Side

Case 1

Server Finished Operation But Client is not available to receive response

Most Servers Discard Response

Case 2

While server is processing log operation, it learns that the client crashed.

1) It previously maintain/track clients status. Once it knows a client crashed, terminates RPC operation
2) Check each RPC process with time intervals. If one takes so long, it checks on client, and delete operations if client crashed
3) Server broadcasts to check if client joined with a different process ID
4) Server sets a timeout at the start, and checks on the client at that time, if it is still alive, continue computation with a new timeout. Repeat until done

Source: https://lass.cs.umass.edu/~shenoy/courses/spring13/lectures/notes/677_lec09.pdf
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RPC Failures – Server Side

Server Could crash before or after work is done

We studied 3 semantics to determine how frequently the remote operation can run in the event of failure

Photos Credit: https://www.datanami.com/2017/02/09/exactly-big-deal-apache-kafka/
RPC Server Failure Semantics *at-least-once*

- If the **client received a reply** from the server, it means that **call has been executed at least once** on the server.
- Doesn’t apply **duplicate request filtering**
  - Whenever **client times out**, it will **resend request**
- If **server crashed** during execution or after execution but before sending reply
  - Server **comes up and then again executes** the call and **sends the reply** to the client.
  - That is why it is Suitable For **Idempotent** Operations
- How does it **make sure client received the reply**?
  - It waits for client **ACK**. If **lost**, and server **times out**, it will **retransmit reply**

- **Use Cases**: Applications where receiving every message is important & duplicity is acceptable (Queries)

**Worked**

**Didn’t Work**

Carnegie Mellon University Qatar
**RPC Server Failure Semantics** *at-most-once*

- If the **client gets a reply** from the server, it means that the RPC call have been **executed at most once** on the server.
- If **timeout doesn’t happen** at client, it means server must have **succeeded** in the first try.
- Applies **duplicate requests filtering**
- Suitable for **non-idempotent operations**
- When **operation request is received**
  - If **executed before**, it looks it up from non-volatile memory & sends it
  - If **not executed before**, process it
- If **server crashes** during execution?
  - If a failure happens before all actions are executed, roll back (ensures **atomicity**)
  - If a failure happens while results are in volatile memory, Redo (ensures **durability**)
- **It doesn’t worry about client receiving the reply**
- **Use Cases**: Applications that need high throughput and low latency (collecting sensor measurements)

**Worked**

**Didn’t Work (timeout)**
RPC Server Failure Semantics *exactly-once*

- This is the **most desirable** scenario, no matter what happened on the server side (crashes / restarts)
- If the client gets back a reply, then it is confirmed that the RPC call was **executed exactly once** on the server
- Applies **duplicate requests filtering**
- Server **keeps results** in cache **until client ACKs**
- Most **like the local-procedure calls** BUT **very expensive** to implement!
  - That’s why systems implement weaker semantics and errors need to be handled explicitly
- Use Cases: **Financial Operations**
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**Bandwidth v.s. Throughput**

**Bandwidth**: The maximum speed of bit transmission in a channel or a link, measured in **bits/sec**.

**Throughput**: The actual speed of bit transmission in a channel or a link, measured in **bits/sec**.

Why are they different?

- Because networks often experience issues that hinder performance (e.g., overloading).
RPC Time - Communication Time

**Transmission time:** The time it takes to upload your message from your machine to the transferring medium (link)
- Measured from when first bit of data is pushed onto the link until last bit of data is pushed onto the link

\[ \text{Transmission time} = \frac{\text{Message Size}}{\text{Throughput}} \]
(P.S: Throughput might be equal to Bandwidth in some situations)

**Propagation time:** The time it takes for a bit in your message to travel across the transferring medium from source to destination.

\[ \text{Propagation time} = \frac{\text{Distance}}{\text{Speed of Light}} \]

**Communication time** = **Transmission time** + **Propagation time**
RPC Total Time

\[
\text{RPC Total Time} = \text{TimeToSendRequest} + \text{TimeToReceiveRequest} + \text{TimeToProcess} + \text{TimeToSendResults} + \text{TimeToReceiveResult}
\]
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<table>
<thead>
<tr>
<th>Propagation delay depends on ___</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bandwidth</td>
</tr>
<tr>
<td>Packet length</td>
</tr>
<tr>
<td>Distance between the routers</td>
</tr>
<tr>
<td>CPU speed</td>
</tr>
<tr>
<td>None of the above</td>
</tr>
</tbody>
</table>
Transmission delay does not depend on ____________

- Packet length
- Distance between the routers
- Bandwidth / Throughput of medium
- None of the above
As Bandwidth $\times$ Length increases, uncertainty increases
A network with bandwidth of 10 Mbps can pass 12,000 frames per minute with each frame carrying 10,000 bits. What is the throughput of this network as a percentage of the bandwidth?

10%
12%
20%
25%
Calculate the propagation time needed for a bit to travel between two points with a connecting cable length of 12,000 m. Assuming that the cable's propagation speed is $2.4 \times 10^8$ m/s.

- 50 ms
- 5 ms
- 0.5 ms
- 0.05 ms
Sender "ping"s a receiver in a different country. Ping would send a packet of size 64 bytes. What is the minimum time that it will take the ping packet to return to the sender?

<table>
<thead>
<tr>
<th>Link</th>
<th>One-way propagation delay</th>
<th>Capacity</th>
<th>Time Calculation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Link A</td>
<td>10ms</td>
<td>1Gbit/second</td>
<td>((2 \times 64) / 1073741824 + 30 \text{ ms})</td>
</tr>
<tr>
<td>Link B</td>
<td>20ms</td>
<td>1Gbit/second</td>
<td>((2 \times 64 \times 8) / 1073741824 + 30 \text{ ms})</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>((4 \times 64) / 1073741824 + 60 \text{ ms})</td>
</tr>
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</tr>
</tbody>
</table>
Which RPC semantic would you choose for the below situations: Transferring money to your friend?

- Exactly-once
- At-least-once
- At-most-once
- None
Which RPC semantic would you choose for the below situations: Posting a message on messenger?

- At-least-once
- At-most-once
- Exactly-once
- Nonce
Which RPC semantic would you choose for the below situations: Deleting a friend from social media?

- At-least-once
- At-most-once
- Exactly-once
- Nonce
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Your Feedback & Suggestions

- More Office Hours, none of the OH suit my schedule
- Maybe clarify some of the questions
- Just a bit more detail
- Go through theory a bit faster so that we can get to understanding the implementation more. As typically with theory everything makes sense but when it comes to implementation, I struggle a lot more, especially if it something unfamiliar.
- Not sure is this is a good idea, but it could possibly be useful to have a handout/very short questionnaire about the content of the recitation to check our understanding.
Reflection & Plan

• Office Hours & Availability
  • Added More Office Hours (Open Door): every day!
    • Sunday & Thursday : 10-12,
    • Tuesday 12-1
    • Monday & Wednesday 9-10
  • Besides, I am daily in the office 7:30-3:30
    • If door is slightly closed (it helps me focus):
      • Knock OR Send me on Google Chat
      • I will meet if time permits
    • Otherwise, we can always schedule a meeting
  • I am always available to answer questions on Piazza
Reflection & Plan

• Suggestion: More on the implementation than theory
  • Depending on the assignment we do some implementation (e.g. Multi-threaded TCP, RMI, ...)
  • During recitations, please point out where more implementation information is needed, I can accommodate if possible
  • We must leave some assignment implementation details for you to figure out!

• Suggestions: Questionnaire to check understanding of concepts at the end of each recitation
  • Will make sure always to include some exercises

• Grading & Feedback
  • Will give more feedback on grades
  • Always happy to explain missed questions
  • Always happy to reassess
The presentations were very clear, and the examples given were very helpful to understanding the concepts needed in project 1.

Going over a topic until everyone understands

clear instructions on approaching the projects, some implementation starters are provided during recitations

Understanding the theory and then doing a practical example and not skipping what ‘simple’ functions do as not everyone has full understanding what every function does.

I like how we go into specific details and the code during recitations.

Useful recitations with filtered information that is useful

contents are organized and explained well

Yes, the OH and feedback is very useful.

By constantly asking if anyone has any questions after each topic

Giving time to ask questions

TA seems like they care about students and makes sure they entirely understand everything.

Small class size and calm manner of explaining (creates an atmosphere where questions are welcomed)
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