

# Introduction to Computer Systems

15-213, fall 2009

7<sup>th</sup> Lecture, Sep. 14<sup>th</sup>

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# Last Time

## ■ For loops

- for loop → while loop → do-while loop → goto version
- for loop → while loop → goto “jump to middle” version

## ■ Switch statements

- Jump tables: `jmp *.L62(,%edx,4)`
- Decision trees (not shown)

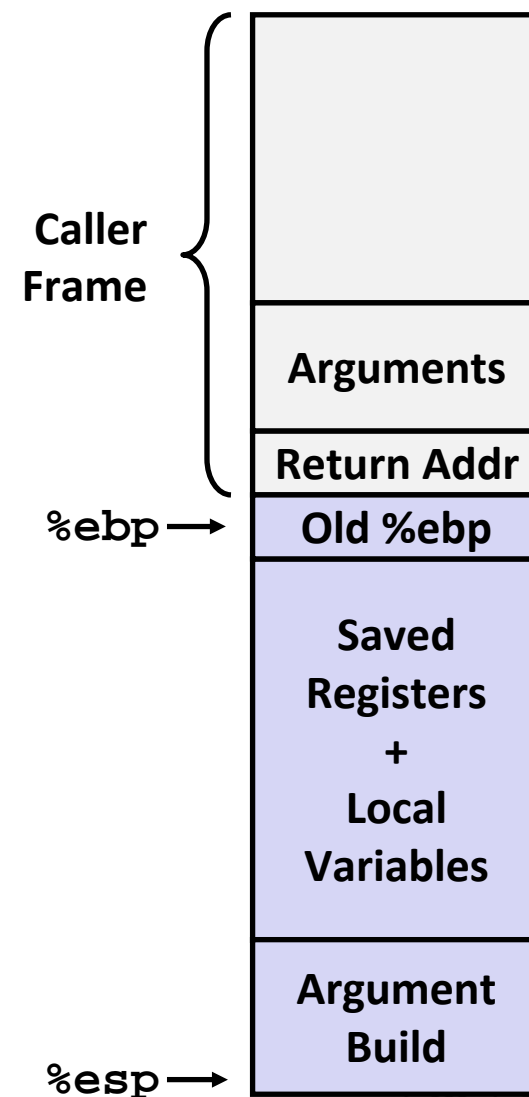
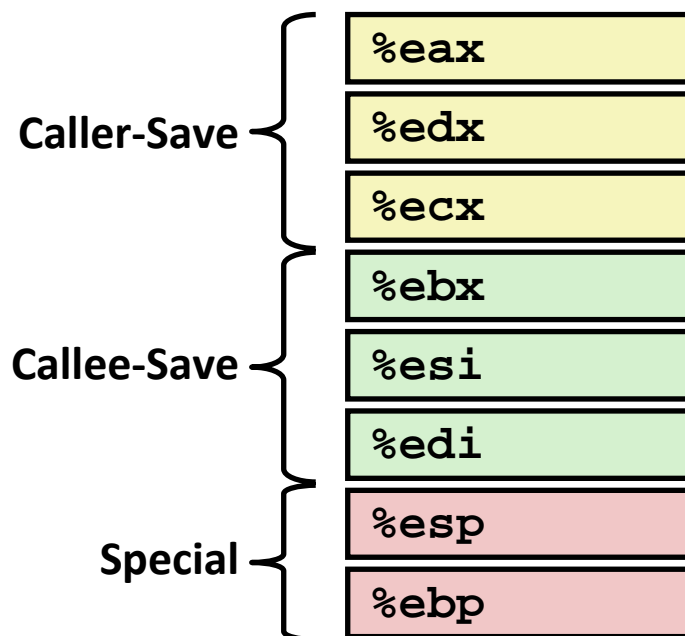
### Jump table

```
.section .rodata
    .align 4
.L62:
    .long    .L61    # x = 0
    .long    .L56    # x = 1
    .long    .L57    # x = 2
    .long    .L58    # x = 3
    .long    .L61    # x = 4
    .long    .L60    # x = 5
    .long    .L60    # x = 6
```

# Last Time

## ■ Procedures (IA32)

- call / return
- %esp, %ebp
- local variables
- recursive functions



# Today

## ■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

## ■ Structures

# Basic Data Types

## ■ Integral

- Stored & operated on in general (integer) registers
- Signed vs. unsigned depends on instructions used

Intel	GAS	Bytes	C
byte	<b>b</b>	1	<b>[unsigned] char</b>
word	<b>w</b>	2	<b>[unsigned] short</b>
double word	<b>l</b>	4	<b>[unsigned] int</b>
quad word	<b>q</b>	8	<b>[unsigned] long int (x86-64)</b>

## ■ Floating Point

- Stored & operated on in floating point registers

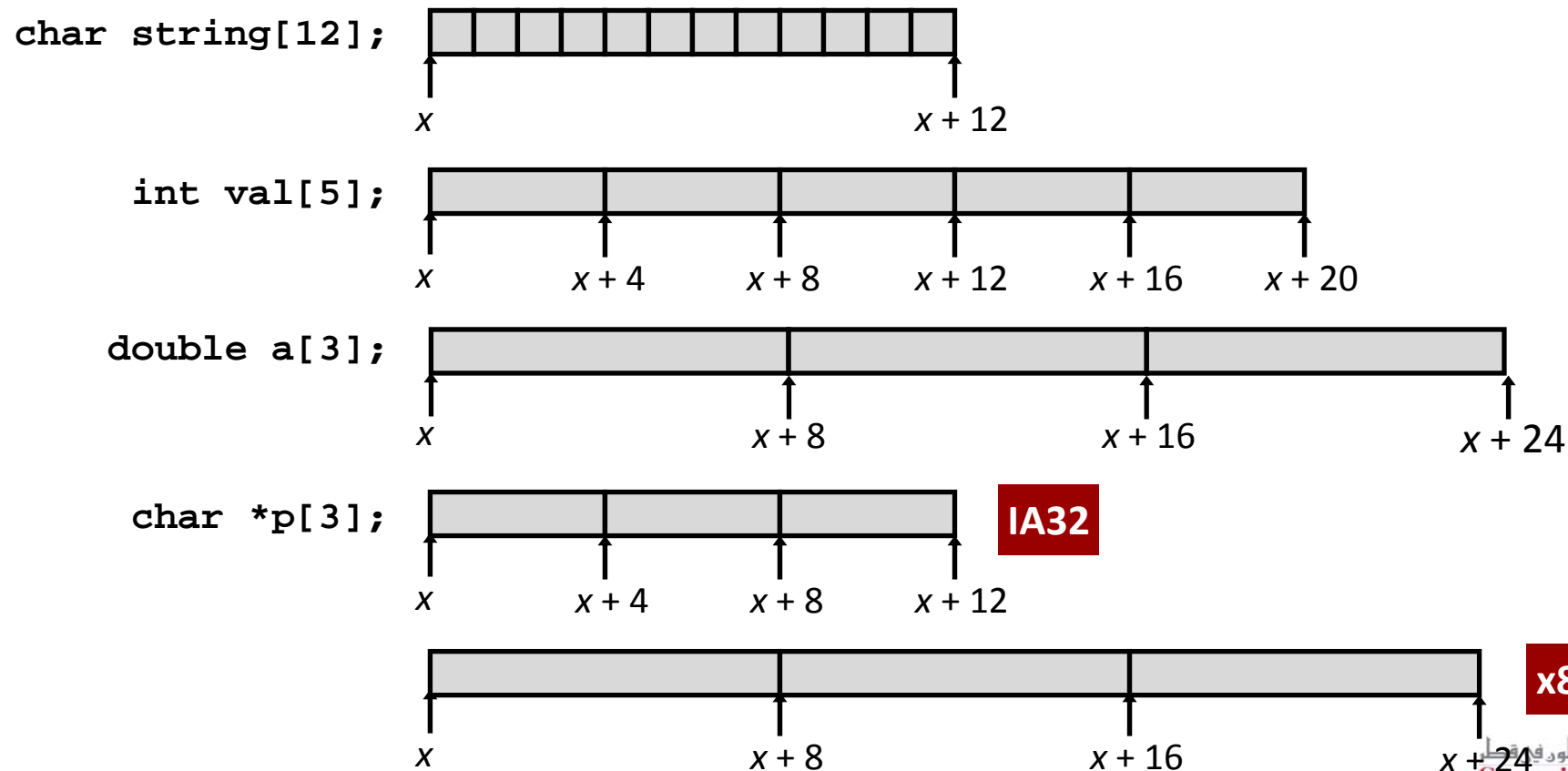
Intel	GAS	Bytes	C
Single	<b>s</b>	4	<b>float</b>
Double	<b>l</b>	8	<b>double</b>
Extended	<b>t</b>	10/12/16	<b>long double</b>

# Array Allocation

## ■ Basic Principle

$T$   $A[L];$

- Array of data type  $T$  and length  $L$
- Contiguously allocated region of  $L * \text{sizeof}(T)$  bytes

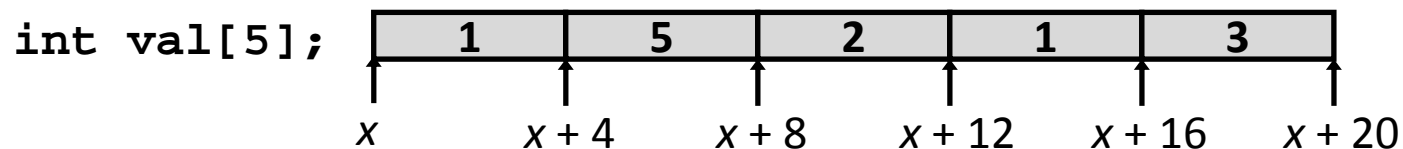


# Array Access

## ■ Basic Principle

$T$   $A[L];$

- Array of data type  $T$  and length  $L$
- Identifier  $A$  can be used as a pointer to array element 0: Type  $T^*$



## ■ Reference

Type

Value

`val[4]`

`val`

`val+1`

`&val[2]`

`val[5]`

`*(val+1)`

`val + i`

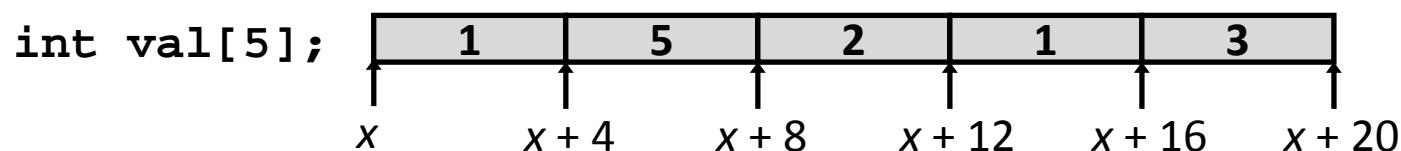
Will disappear  
Blackboard?

# Array Access

## ■ Basic Principle

$T$   $A[L];$

- Array of data type  $T$  and length  $L$
- Identifier  $A$  can be used as a pointer to array element 0: Type  $T^*$



## ■ Reference      Type      Value

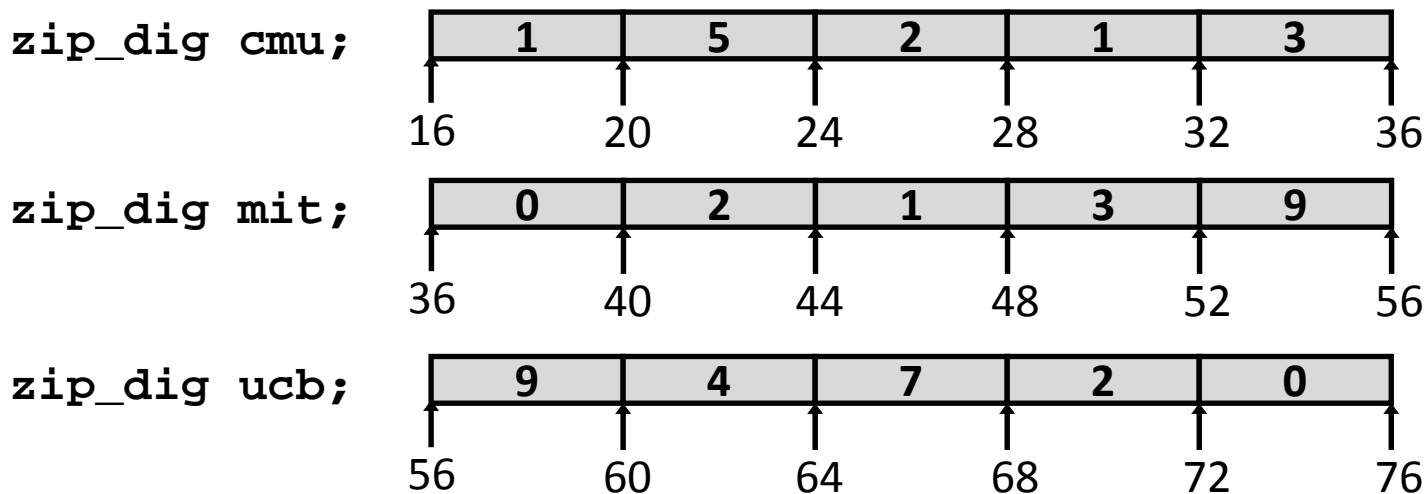
<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	$x$
<code>val+1</code>	<code>int *</code>	$x+4$
<code>&amp;val[2]</code>	<code>int *</code>	$x+8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val + i</code>	<code>int *</code>	$x+4i$



# Array Example

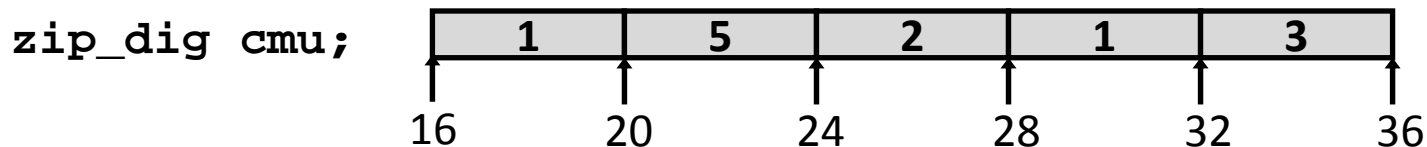
```
typedef int zip_dig[5];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “zip\_dig cmu” equivalent to “int cmu[5]”
- Example arrays were allocated in successive 20 byte blocks
  - Not guaranteed to happen in general

# Array Accessing Example



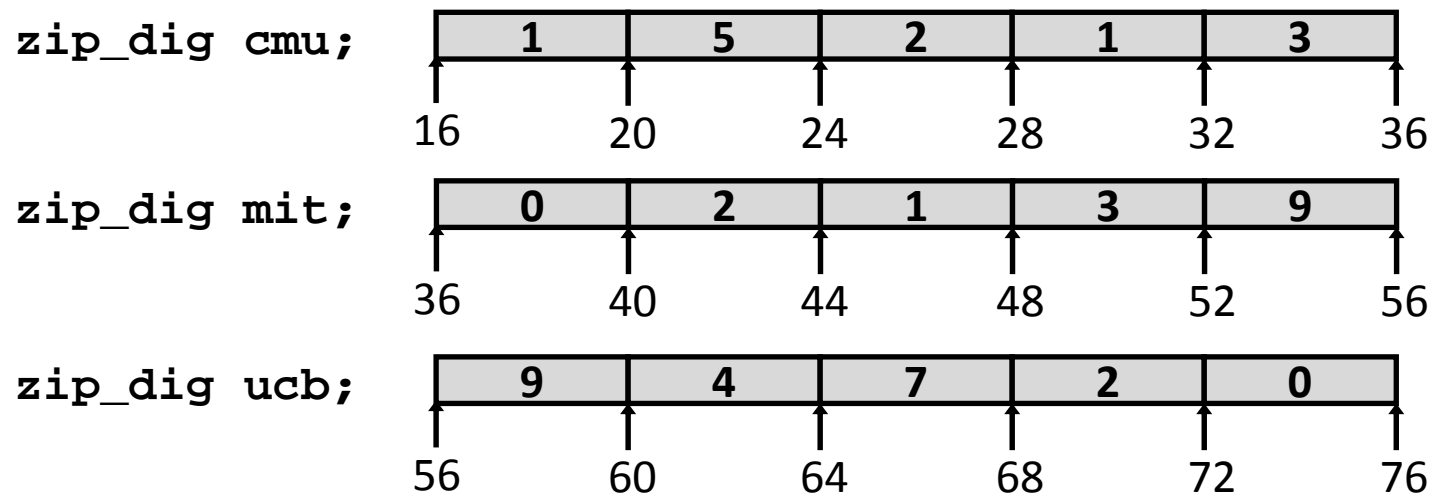
```
int get_digit
  (zip_dig z, int dig)
{
  return z[dig];
}
```

## IA32

```
# %edx = z
# %eax = dig
movl (%edx,%eax,4),%eax #
z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at  $4 * \%eax + \%edx$
- Use memory reference  $(\%edx, \%eax, 4)$

# Referencing Examples



## Reference

mit[3]  
mit[5]  
mit[-1]  
cmu[15]

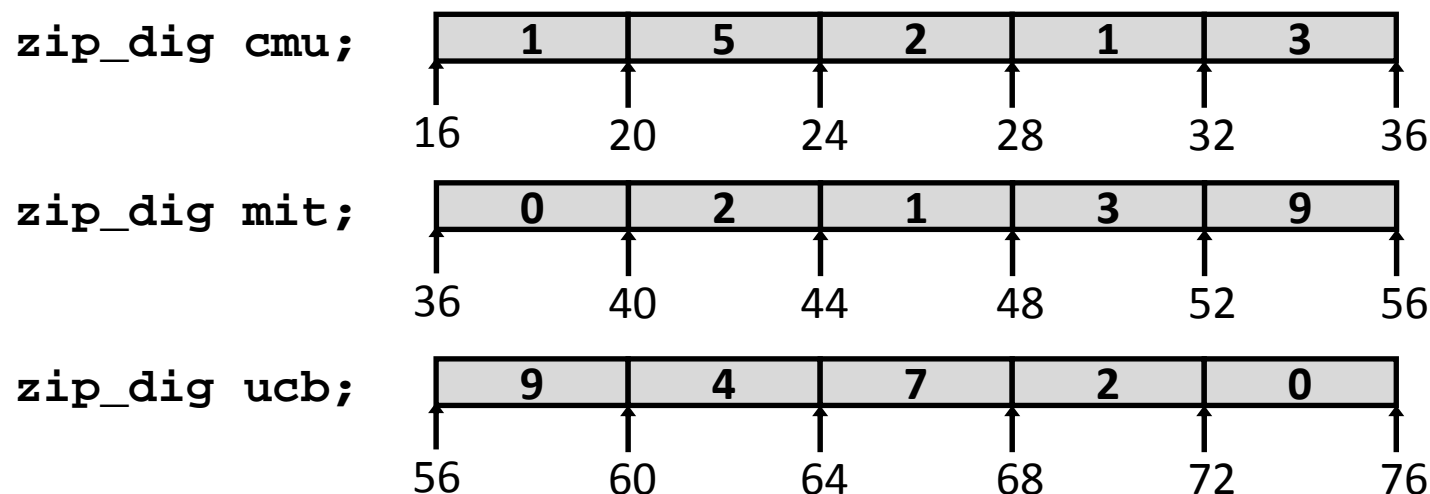
Address

Value

Guaranteed?

Will disappear  
Blackboard?

# Referencing Examples



Reference	Address	Value	Guaranteed?
mit[3]	$36 + 4 * 3 = 48$	3	Yes
mit[5]	$36 + 4 * 5 = 56$	9	No
mit[-1]	$36 + 4 * -1 = 32$	3	No
cmu[15]	$16 + 4 * 15 = 76$	??	No

- No bound checking
- Out of range behavior implementation-dependent
- No guaranteed relative allocation of different arrays

# Array Loop Example

## ■ Original

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

## ■ Transformed

- As generated by GCC
- Eliminate loop variable *i*
- Convert array code to pointer code
- Express in do-while form (no test at entrance)

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while (z <= zend);
    return zi;
}
```

# Array Loop Implementation (IA32)

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

```
# %ecx = z
xorl %eax,%eax
leal 16(%ecx),%ebx
.L59:
leal (%eax,%eax,4),%edx
movl (%ecx),%eax
addl $4,%ecx
leal (%eax,%edx,2),%eax
cmpl %ebx,%ecx
jle .L59
```

**Will disappear  
Blackboard?**

# Array Loop Implementation (IA32)

## ■ Registers

`%ecx` `z`

`%eax` `zi`

`%ebx` `zend`

## ■ Computations

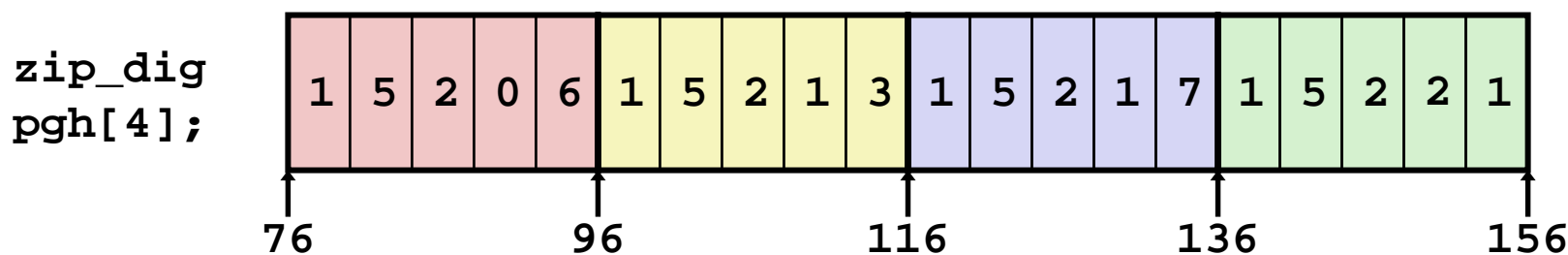
- $10 * z_i + *z$  implemented as  $*z + 2 * (z_i + 4 * z_i)$
- `z++` increments by 4

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

```
# %ecx = z
xorl %eax,%eax           # zi = 0
leal 16(%ecx),%ebx       # zend = z+4
.L59:
leal (%eax,%eax,4),%edx  # 5*zi
movl (%ecx),%eax         # *z
addl $4,%ecx             # z++
leal (%eax,%edx,2),%eax  # zi = *z + 2*(5*zi)
cmpl %ebx,%ecx         # z : zend
jle .L59              # if <= goto loop
```

# Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3},
   {1, 5, 2, 1, 7},
   {1, 5, 2, 2, 1}};
```



- “zip\_dig pgh[4]” equivalent to “int pgh[4][5]”
  - Variable `pgh`: array of 4 elements, allocated contiguously
  - Each element is an array of 5 `int`'s, allocated contiguously
- “Row-Major” ordering of all elements guaranteed



# Multidimensional (Nested) Arrays

## ■ Declaration

$T$   $A[R][C];$

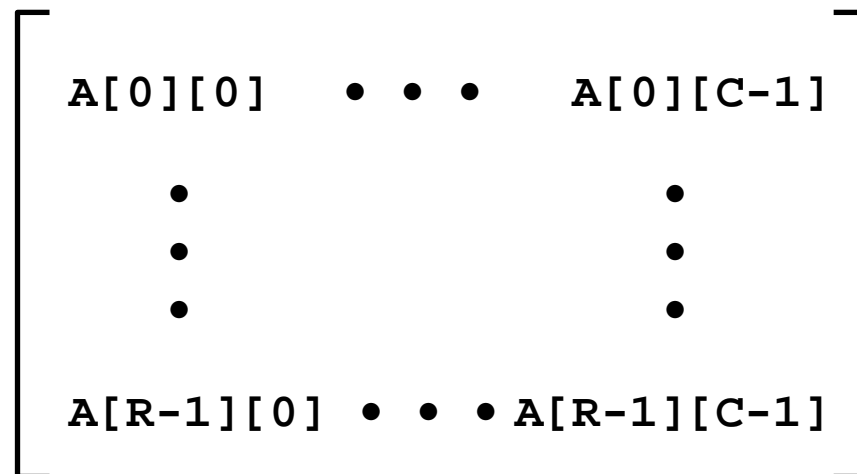
- 2D array of data type  $T$
- $R$  rows,  $C$  columns
- Type  $T$  element requires  $K$  bytes

## ■ Array Size

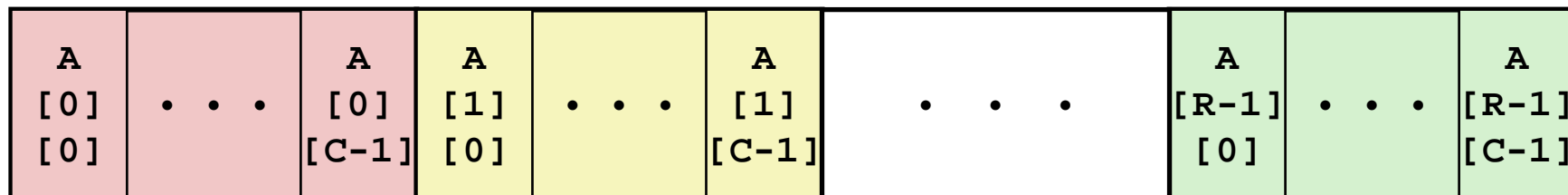
- $R * C * K$  bytes

## ■ Arrangement

- Row-Major Ordering



`int A[R][C];`



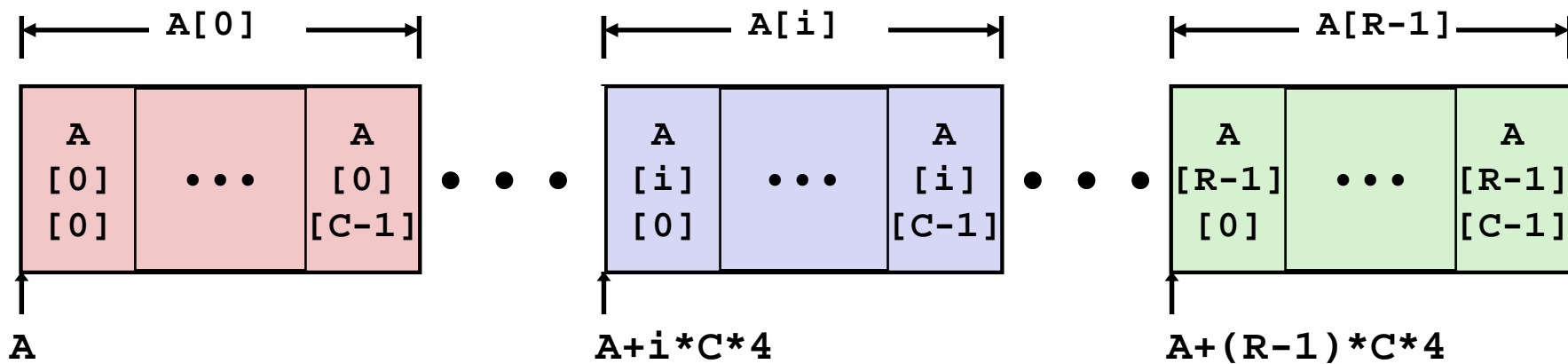
$4 * R * C$  Bytes

# Nested Array Row Access

## ■ Row Vectors

- $A[i]$  is array of  $C$  elements
- Each element of type  $T$  requires  $K$  bytes
- Starting address  $A + i * (C * K)$

```
int A[R][C];
```



# Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 }},
     {1, 5, 2, 1, 7 }},
     {1, 5, 2, 2, 1 }};
```

- What data type is `pgh[index]`?
- What is its starting address?

```
# %eax = index
leal (%eax,%eax,4),%eax
leal pgh(,%eax,4),%eax
```

Will disappear  
Blackboard?

# Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
    {{1, 5, 2, 0, 6},
     {1, 5, 2, 1, 3 },
     {1, 5, 2, 1, 7 },
     {1, 5, 2, 2, 1 }};
```

```
# %eax = index
leal (%eax,%eax,4),%eax # 5 * index
leal pgh(,%eax,4),%eax # pgh + (20 * index)
```

## ■ Row Vector

- `pgh[index]` is array of 5 `int`'s
- Starting address `pgh+20*index`

## ■ IA32 Code

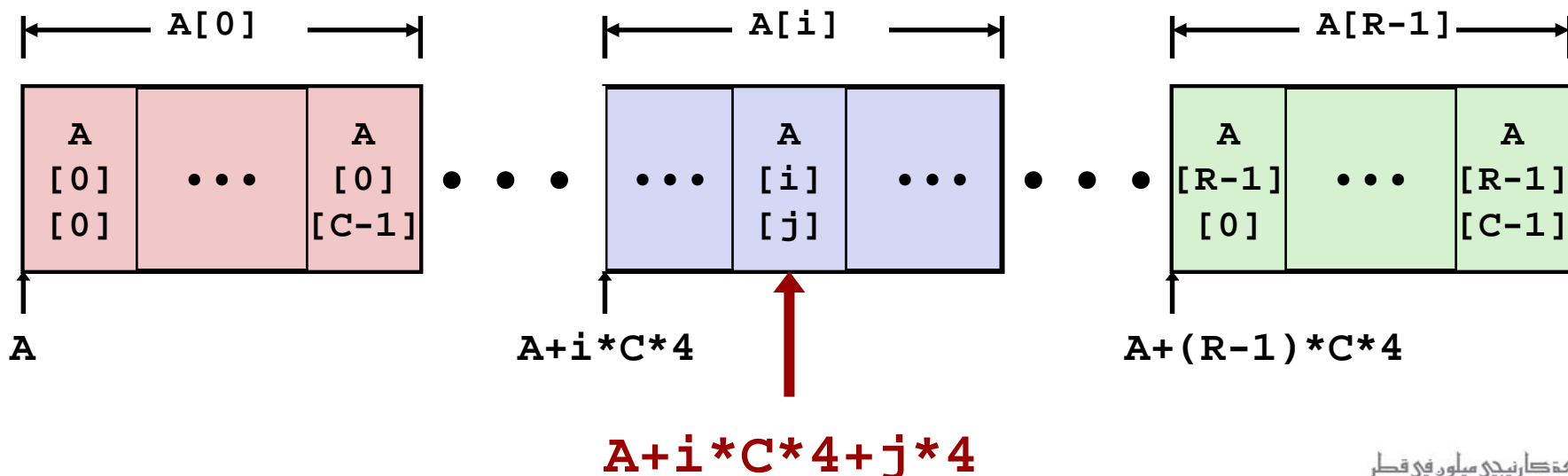
- Computes and returns address
- Compute as `pgh + 4*(index+4*index)`

# Nested Array Row Access

## ■ Array Elements

- $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
- Address  $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



# Nested Array Element Access Code

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```

```
# %ecx = dig
# %eax = index
leal 0(,%ecx,4),%edx      # 4*dig
leal (%eax,%eax,4),%eax   # 5*index
movl pgh(%edx,%eax,4),%eax # *(pgh + 4*dig + 20*index)
```

## ■ Array Elements

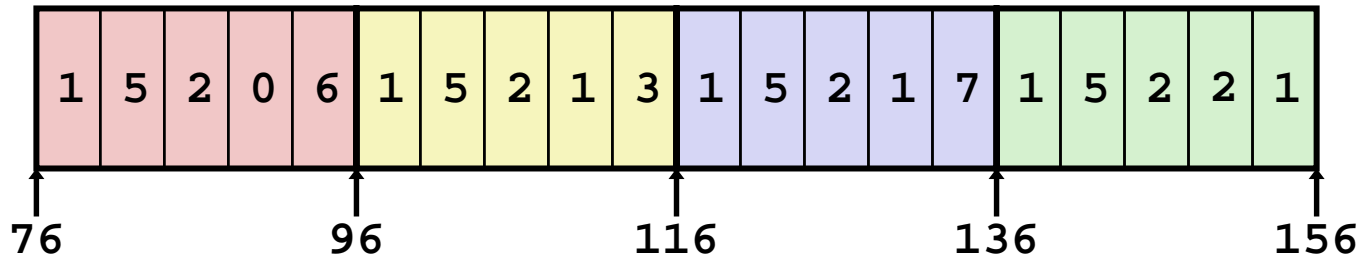
- `pgh[index][dig]` is `int`
- Address: `pgh + 20*index + 4*dig`

## ■ IA32 Code

- Computes address `pgh + 4*dig + 4*(index+4*index)`
- `movl` performs memory reference

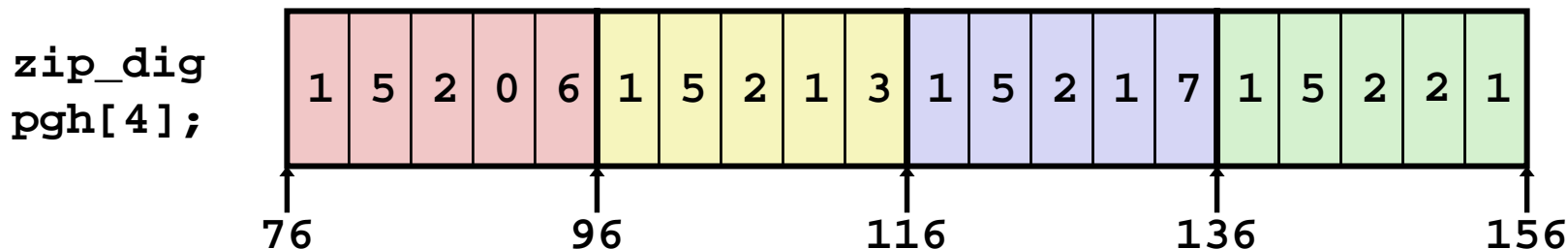
# Strange Referencing Examples

```
zip_dig
pgh[4];
```



Reference	Address	Value	Guaranteed?
<code>pgh[3][3]</code>	<b>Will disappear</b>		
<code>pgh[2][5]</code>			
<code>pgh[2][-1]</code>			
<code>pgh[4][-1]</code>			
<code>pgh[0][19]</code>			
<code>pgh[0][-1]</code>			

# Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>pgh[3][3]</code>	$76+20*3+4*3 = 148$	2	Yes
<code>pgh[2][5]</code>	$76+20*2+4*5 = 136$	1	Yes
<code>pgh[2][-1]</code>	$76+20*2+4*-1 = 112$	3	Yes
<code>pgh[4][-1]</code>	$76+20*4+4*-1 = 152$	1	Yes
<code>pgh[0][19]</code>	$76+20*0+4*19 = 152$	1	Yes
<code>pgh[0][-1]</code>	$76+20*0+4*-1 = 72$	??	No

- Code does not do any bounds checking
- Ordering of elements within array guaranteed

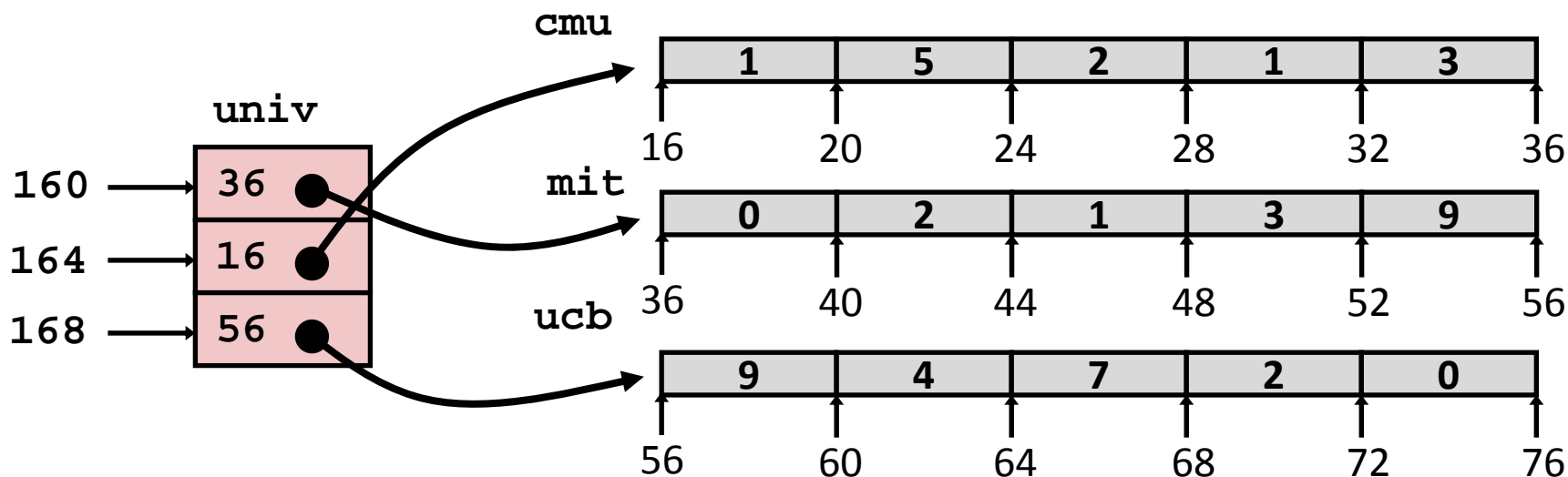


# Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
  - 4 bytes
- Each pointer points to array of `int`'s



# Element Access in Multi-Level Array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

```
# %ecx = index
# %eax = dig
leal 0(,%ecx,4),%edx
movl univ(%edx),%edx
movl (%edx,%eax,4),%eax
```

**Will disappear  
Blackboard?**

# Element Access in Multi-Level Array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

```
# %ecx = index
# %eax = dig
leal 0(,%ecx,4),%edx      # 4*index
movl univ(%edx),%edx      # Mem[univ+4*index]
movl (%edx,%eax,4),%eax   # Mem[...+4*dig]
```

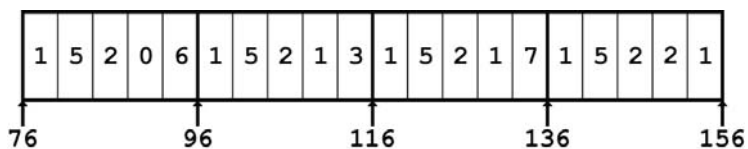
## ■ Computation (IA32)

- Element access `Mem[Mem[univ+4*index]+4*dig]`
- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

# Array Element Accesses

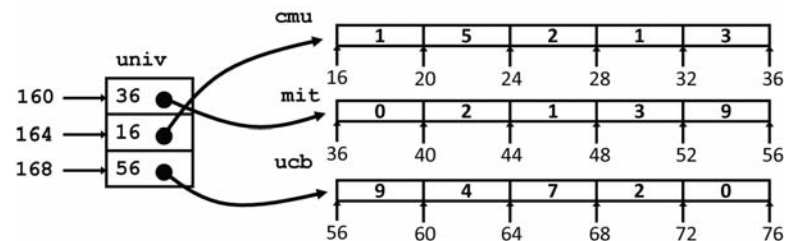
## Nested array

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```



## Multi-level array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

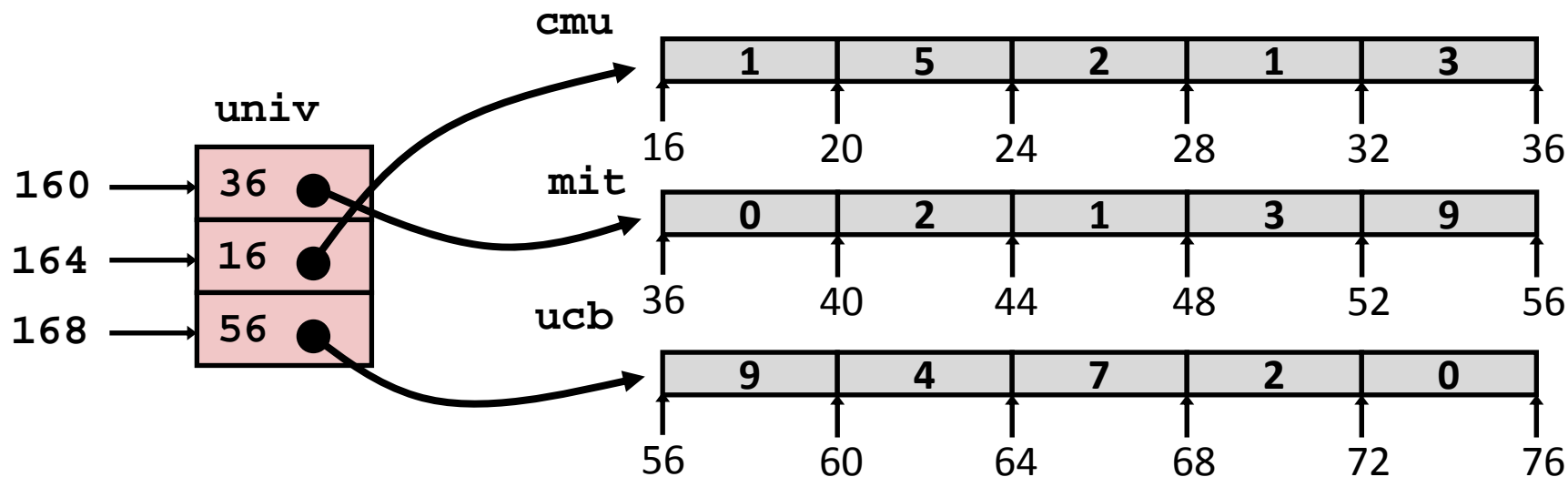


Access looks similar, but element:

$\text{Mem}[\text{pgh} + 20 * \text{index} + 4 * \text{dig}]$

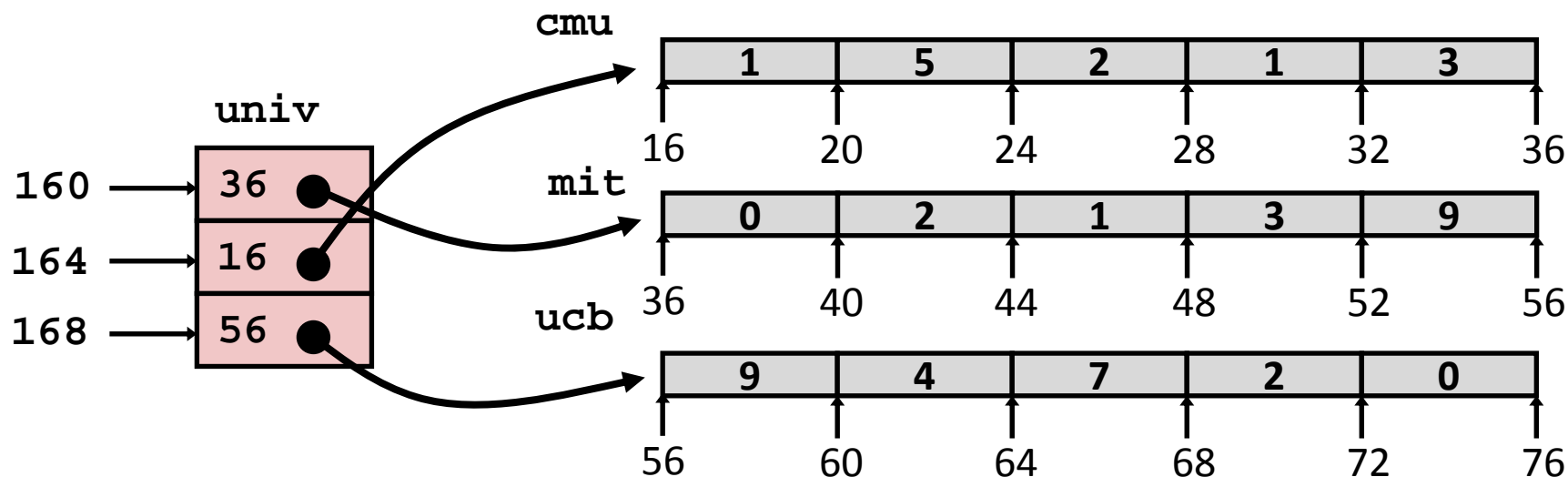
$\text{Mem}[\text{Mem}[\text{univ} + 4 * \text{index}] + 4 * \text{dig}]$

# Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>univ[2][3]</code>	<b>Will disappear</b>		
<code>univ[1][5]</code>			
<code>univ[2][-1]</code>			
<code>univ[3][-1]</code>			
<code>univ[1][12]</code>			

# Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>univ[2][3]</code>	$56+4*3 = 68$	2	Yes
<code>univ[1][5]</code>	$16+4*5 = 36$	0	No
<code>univ[2][-1]</code>	$56+4*-1 = 52$	9	No
<code>univ[3][-1]</code>	??	??	No
<code>univ[1][12]</code>	$16+4*12 = 64$	7	No

- Code does not do any bounds checking
- Ordering of elements in different arrays not guaranteed

# Using Nested Arrays

## ■ Strengths

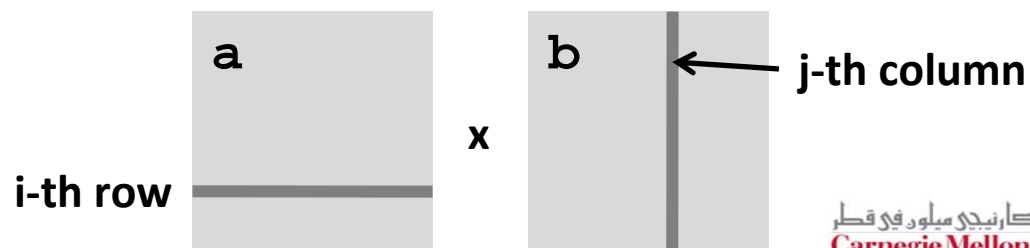
- C compiler handles doubly subscripted arrays
- Generates very efficient code
- Avoids multiply in index computation

## ■ Limitation

- Only works for fixed array size

```
#define N 16
typedef int fix_matrix[N][N];
```

```
/* Compute element i,k of
   fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
 int i, int k)
{
    int j;
    int result = 0;
    for (j = 0; j < N; j++)
        result += a[i][j]*b[j][k];
    return result;
}
```



# Dynamic Nested Arrays

## ■ Strength

- Can create matrix of any size

## ■ Programming

- Must do index computation explicitly

## ■ Performance

- Accessing single element costly
- Must do multiplication

```
int * new_var_matrix(int n)
{
    return (int *)
        calloc(sizeof(int), n*n);
}
```

```
int var_ele
(int *a, int i, int j, int n)
{
    return a[i*n+j];
}
```

```
movl 12(%ebp),%eax    # i
movl 8(%ebp),%edx     # a
imull 20(%ebp),%eax   # n*i
addl 16(%ebp),%eax    # n*i+j
movl (%edx,%eax,4),%eax # Mem[a+4*(i*n+j)]
```



# Dynamic Array Multiplication

## ■ Without Optimizations

- Multiplies: 3
  - 2 for subscripts
  - 1 for data
- Adds: 4
  - 2 for array indexing
  - 1 for loop index
  - 1 for data

```
/* Compute element i,k of
   variable matrix product */
int var_prod_ele
(int *a, int *b,
 int i, int k, int n)
{
    int j;
    int result = 0;
    for (j = 0; j < n; j++)
        result +=
            a[i*n+j] * b[j*n+k];
    return result;
}
```

# Optimizing Dynamic Array Multiplication

## ■ Optimizations

- Performed when set optimization level to `-O2`

## ■ Code Motion

- Expression `i*n` can be computed outside loop

## ■ Strength Reduction

- Incrementing `j` has effect of incrementing `j*n+k` by `n`

## ■ Operations count

- 4 adds, 1 mult

## ■ Compiler can optimize regular access patterns

```
{
    int j;
    int result = 0;
    for (j = 0; j < n; j++)
        result +=
            a[i*n+j] * b[j*n+k];
    return result;
}
```

```
{
    int j;
    int result = 0;
    int iTn = i*n;
    int jTnPk = k;
    for (j = 0; j < n; j++) {
        result +=
            a[iTn+j] * b[jTnPk];
        jTnPk += n;
    }
    return result;
}
```

# Today

## ■ Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

## ■ Structures

# Structures

```
struct rec {
    int i;
    int a[3];
    int *p;
};
```

## Memory Layout



## ■ Concept

- Contiguously-allocated region of memory
- Refer to members within structure by names
- Members may be of different types

## ■ Accessing Structure Member

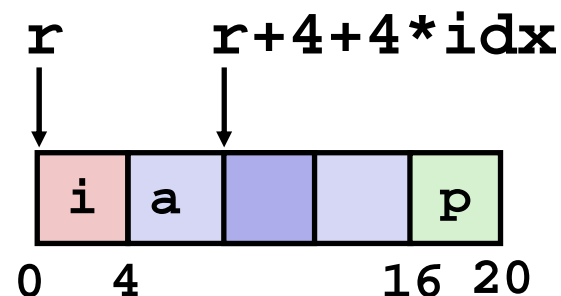
```
void
set_i(struct rec *r,
      int val)
{
    r->i = val;
}
```

## IA32 Assembly

```
# %eax = val
# %edx = r
movl %eax, (%edx)    # Mem[r] = val
```

# Generating Pointer to Structure Member

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```



## ■ Generating Pointer to Array Element

- Offset of each structure member determined at compile time

```
int *find_a
(struct rec *r, int idx)
{
  return &r->a[idx];
}
```

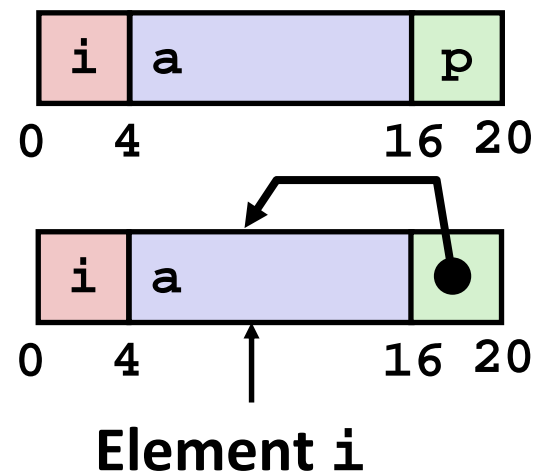
```
# %ecx = idx
# %edx = r
leal 0(,%ecx,4),%eax # 4*idx
leal 4(%eax,%edx),%eax # r+4*idx+4
```

# Structure Referencing (Cont.)

## ■ C Code

```
struct rec {
    int i;
    int a[3];
    int *p;
};
```

```
void
set_p(struct rec *r)
{
    r->p =
        &r->a[r->i];
}
```



```
# %edx = r
movl (%edx),%ecx      # r->i
leal 0(,%ecx,4),%eax  # 4*(r->i)
leal 4(%edx,%eax),%eax # r+4+4*(r->i)
movl %eax,16(%edx)   # Update r->p
```