

Visualizing heaps

Use the visualization at <http://www.cs.usfca.edu/~galles/visualization/Heap.html> to insert the following elements into a min-heap, in the given order.

5, 3, 6, 7, 2, 6

Priority queue client and library interface

We use heaps to efficiently implement the *priority queue* interface.

```

/* Client interface */
// typedef _____ elem;
typedef void* elem;

// f(x,y) returns true if e1 has STRICTLY higher priority than e2
typedef bool has_higher_priority_fn(elem e1, elem e2);

/* Library interface */
//typedef _____* pq_t;

bool pq_empty(pq_t H);
bool pq_full(pq_t H);
pq_t pq_new(int capacity, has_higher_priority_fn* priority)
    /*@requires capacity > 0 && priority != NULL; @*/
    /*@ensures pq_empty(\result); @*/ ;

void pq_add(pq_t H, elem x)
    /*@requires !pq_full(H); @*/ ;

elem pq_rem(pq_t H) // Removes highest priority element
    /*@requires !pq_empty(H); @*/

```

Checkpoint 0

If the client's `elem` type is picked to be `void*`, will this client interface cause `pq_new(20, &higher_priority)` to return a min-heap, a max-heap, or something else?

```

1 bool higher_priority(void* x, void* y)
2 /*@requires x != NULL && \hastag(int*, x);
3 /*@requires y != NULL && \hastag(int*, y);
4 {
5     return *(int*)x > *(int*)y;
6 }

```

Checkpoint 1

Write a function (of type `has_higher_priority_fn`) that ensures that, in a priority queue of (pointers to) strings, the *longest* strings always gets returned first. The `string_length` function may be helpful.

Deletion of the highest-priority element from a heap

You may need the following functions: `is_heap_safe`, `ok_above`

```
1 void sift_down(heap* H)
2 //@requires _____;
3 //@ensures is_heap(H);
4 {
5     int i = 1;
6     while ( _____ )
7         //@loop_invariant 1 <= i && i < H->next;
8         //@loop_invariant is_heap_except_down(H, i);
9         //@loop_invariant grandparent_check(H, i);
10    {
11        int left = 2*i;
12        int right = left+1;
13        if ( _____ )
14            return;
15        if ( _____ ) {
16            swap_up(H, left);
17            i = left;
18        } else {
19            //@assert _____;
20            swap_up(H, right);
21            i = right;
22        }
23    }
24 }
25
26 elem pq_rem(heap* H)
27 //@requires is_heap(H) && !pq_empty(H);
28 //@ensures is_heap(H);
29 {
30     elem min = H->data[1];
31     (H->next)--;
32     if (H->next > 1) {
33         H->data[1] = H->data[H->next];
34         sift_down(H);
35     }
36     return min;
37 }
```

Checkpoint 2

- (a) Check that the preconditions imply the loop invariants hold initially, and that they are satisfied when `sift_down` is called from `pq_rem`.
- (b) Show that the grandparent check is necessary as a loop invariant.
- (c) Prove that the loop invariants imply the postcondition for the return on line 14 and on line 23.